

GAMES CHARACTERIZING LIMSUP FUNCTIONS AND BAIRE CLASS 1 FUNCTIONS

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Abstract. We consider a real-valued function f defined on the set of infinite branches X of a countably branching pruned tree T . The function f is said to be a *limsup function* if there is a function $u: T \rightarrow \mathbb{R}$ such that $f(x) = \limsup_{t \rightarrow \infty} u(x_0, \dots, x_t)$ for each $x \in X$. We study a game characterization of limsup functions, as well as a novel game characterization of functions of Baire class 1.

§1. Introduction. Throughout the paper, let T be a pruned tree on a non-empty countable set A , and X be the set of its infinite branches. We say that $f: X \rightarrow \mathbb{R}$ is a *limsup function* if there exists a function $u: T \rightarrow \mathbb{R}$ such that, for every $x \in X$,

$$f(x) = \limsup_{t \rightarrow \infty} u(x_0, \dots, x_t). \tag{1.1}$$

Payoff evaluations of limsup type are ubiquitous in gambling theory [3], in the theory of dynamic games [10], and in computer science [1]. Limsup payoff evaluation expresses the decision maker’s preference to receive high payoff infinitely often.

We first relate limsup functions to certain well-known classes of functions. In fact, f is a limsup function precisely if it is a pointwise limit of a descending sequence of lower semicontinuous functions. Pointwise limits of a descending sequence of lower semicontinuous functions have been studied, e.g., in [5]. In particular, it is known that f is a limsup function exactly if its subgraph is a $\mathbf{\Pi}_2^0$ set (i.e., a G_δ set), and that the sum, the minimum, and the maximum of two limsup functions is a limsup function. We also deduce a characterization of Baire class 1 functions $f: X \rightarrow \mathbb{R}$: these are exactly the functions such that both f and $-f$ are limsup functions.

The core of the paper is devoted to the study of two related games. The first one is the following:

I	x_0	x_1	\dots
II	v_0	v_1	\dots

The moves x_0, x_1, \dots of Player I are points of A such that $(x_0, \dots, x_t) \in T$ for each $t \in \mathbb{N}$. The moves v_0, v_1, \dots of Player II are real numbers (the results go through as

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stated or with obvious modifications for a version of the game where Player II is restricted to play rational numbers). The game starts with a move of Player I, x_0 . Having observed x_0 , Player II chooses v_0 . Having observed v_0 , Player I chooses x_1 , and so on. In this fashion the players produce a run of the game, $(x_0, v_0, x_1, v_1, \dots)$. Player II wins the run if $f(x_0, x_1, \dots) = \limsup v_t$. We denote this game by $\Gamma(f)$.

As we will see in Lemma 3.1, Player II has a winning strategy in $\Gamma(f)$ precisely when f is a limsup function. Whether Player I has a winning strategy in $\Gamma(f)$ turns out to be a more subtle question. We give a sufficient condition for Player I to have a winning strategy in $\Gamma(f)$, a condition that is also necessary if either f is Borel measurable (more precisely, it suffices if the sets of the form $\{x \in X : f(x) \geq r\}$ are co-analytic), or if the range of f contains no infinite strictly increasing sequence, in particular if f takes only finitely many values. We also show that the game $\Gamma(f)$ is determined if f is Borel measurable (again, it suffices if the sets of the form $\{x \in X : f(x) \geq r\}$ are co-analytic), but not in general.

The second game, denoted by $\Gamma'(f)$, is as follows:

$$\begin{array}{rcccc} \text{I} & x_0 & & x_1 & & \dots \\ \text{II} & & (v_0, w_0) & & (v_1, w_1) & \dots \end{array}$$

This game is similar to $\Gamma(f)$ except that now the moves $(v_0, w_0), (v_1, w_1), \dots$ of Player II are pairs of real numbers. Player II wins in $\Gamma'(f)$ if $f(x_0, x_1, \dots) = \limsup v_t = \liminf w_t$. We denote this game by $\Gamma'(f)$.

Player II has a winning strategy in the game $\Gamma'(f)$ precisely when he has a winning strategy in both games $\Gamma(f)$ and $\Gamma(-f)$, which is the case exactly when f is in Baire class 1. Moreover, the game $\Gamma'(f)$ is always determined. This result holds for any function f , whether or not f is Borel measurable, and is established without the aid of Martin’s determinacy.

The so-called eraser game characterizing Baire class 1 functions from the Baire space to itself was constructed in [4]. Carroy [2] showed that the eraser game is determined, and Kiss [8] generalized the characterization to functions of arbitrary Polish spaces. Game characterizations of several other classes of functions have been considered in [2, 11, 12].

Section 2 discusses characterizations of limsup functions. Sections 3 and 4 are devoted to the analysis of the games $\Gamma(f)$ and $\Gamma'(f)$, respectively.

Unless stated otherwise, proofs are conducted within ZFC.

§2. Characterizations of limsup functions. For $s \in T$, we let $O(s)$ denote the set of $x \in X$ such that s is an initial segment of x . We refer to $O(s)$ as a cylinder set. We endow X with its usual topology, generated by the base consisting of all cylinder sets. For a function $f : X \rightarrow \mathbb{R}$ write $subgr(f) = \{(x, r) \in X \times \mathbb{R} : f(x) \geq r\}$ to denote the *subgraph* of f . For $r \in \mathbb{R}$, we write $\{f \geq r\} = \{x \in X : f(x) \geq r\}$, $\{f > r\} = \{x \in X : f(x) > r\}$, and $\{f = r\} = \{x \in X : f(x) = r\}$.

THEOREM 2.1. *Consider a function $f : X \rightarrow \mathbb{R}$. The following conditions are equivalent:*

- [C1] *The function f is a limsup function.*
- [C2] *There is a sequence g_0, g_1, \dots of lower semicontinuous functions converging pointwise to f .*

[C3] *There is a non-increasing sequence $g_0 \geq g_1 \geq \dots$ of lower semicontinuous functions converging pointwise to f .*

[C4] *The set $\text{subgr}(f)$ is a Π_2^0 subset of $X \times \mathbb{R}$.*

[C5] *For each $r \in \mathbb{R}$, $\{f \geq r\}$ is a Π_2^0 subset of X .*

We remark that the functions satisfying condition [C5] are sometimes called semi-Borel class 2 (see [7]) or upper semi-Baire class 1 functions. The equivalence of the conditions [C2]–[C5] is in fact well known (see [5]). Below we prove the equivalence of conditions [C1]–[C3].

PROOF THAT [C1] IMPLIES [C2]. Let f be a limsup function, and let u be a function as in (1.1). For $n \in \mathbb{N}$ let $g_n(x) = \sup\{u(x_0, \dots, x_t) : t \geq n\}$. ⊣

PROOF THAT [C2] IMPLIES [C3]. Let g_n be a sequence of lower semicontinuous functions converging pointwise to f . Define $g'_n(x) = \sup\{g_m(x) : m \geq n\}$. This gives a non-increasing sequence of lower semicontinuous functions converging pointwise to f . ⊣

PROOF THAT [C3] IMPLIES [C1]. Consider a non-increasing sequence $g_0 \geq g_1 \geq \dots$ of lower semicontinuous functions converging pointwise to f . We will also assume that for each $n \in \mathbb{N}$, the range of g_n contains only reals of the form $z2^{-n}$ for $z \in \mathbb{Z}$. To see that this could be imposed without loss of generality, consider the functions $g'_n(x) = \min\{z2^{-n} : z \in \mathbb{Z}, g_n(x) \leq z2^{-n}\}$. Then, $\{g'_n > z2^{-n}\}$ is the same as the set $\{g_n > z2^{-n}\}$, implying that g'_n is lower semicontinuous. It is easy to see that $g'_0 \geq g'_1 \geq \dots$ is a non-increasing sequence, and that it converges pointwise to f .

We define the function $u : T \rightarrow \mathbb{R}$. For $n \in \mathbb{N}$ and $r \in \mathbb{R}$ note that the set $\{g_n > r\}$ is an open set, because g_n is assumed to be lower semicontinuous. Take a sequence $s \in T$. Define $R_*(s)$ to be the set of real numbers $r \in \mathbb{R}$ such that $O(s) \subseteq \bigcap_{n \in \mathbb{N}} \{g_n > r\}$. For $n \in \mathbb{N}$ define $R_n(s)$ to be the set of real numbers $r \in \mathbb{R}$ such that $O(s) \subseteq \{g_n > r\}$, and such that for no proper initial segment s' of s does it hold that $O(s') \subseteq \{g_n > r\}$. (We remark that $R_*(s)$ is a half-line and the sets $R_n(s)$ are intervals.) Let $R(s)$ be the union of the sets $R_*(s), R_0(s), R_1(s), \dots$. Notice that the set $R(s)$ is bounded above by $\inf\{g_0(y) : y \in O(s)\}$. If $R(s)$ is non-empty, we define $u(s) = \sup R(s)$. If $R(s)$ is empty, we let $u(s) = -\text{length}(s)$.

We show that u satisfies (1.1). Thus fix an $x \in X$. We write s_t to denote (x_0, \dots, x_t) and let $\alpha = \limsup_{t \rightarrow \infty} u(s_t)$. We must show that $f(x) = \alpha$.

We first show that $f(x) \leq \alpha$.

Take a real number r with $r < f(x)$. We argue that $r \leq \alpha$.

For every $n \in \mathbb{N}$ it holds that $r < g_n(x)$, so $x \in \{g_n > r\}$. Let t_n be the smallest $t \in \mathbb{N}$ for which $O(s_{t_n}) \subseteq \{g_n > r\}$. We distinguish between two cases, depending on whether the sequence t_0, t_1, \dots is bounded or not. Suppose first the sequence t_0, t_1, \dots is unbounded. By the choice of t_n , we have $r \in R_n(s_{t_n})$, and hence $r \leq u(s_{t_n})$. We obtain $r \leq \alpha$, as desired. Suppose now that the sequence t_0, t_1, \dots is bounded, say $t_n \leq t$ for each $n \in \mathbb{N}$. Then, $O(s_t) \subseteq \bigcap_{n \in \mathbb{N}} \{g_n > r\}$. Since for $k \geq t$ the cylinder $O(s_k)$ is contained in $O(s_t)$, we have $r \in R_*(s_k)$, and consequently $r \leq u(s_k)$. We conclude that $r \leq \alpha$, as desired.

We now show that $\alpha \leq f(x)$.

We know that $-\infty < \alpha$. Take a real number $r < \alpha$. We now argue that $r \leq f(x)$.

There exists an increasing sequence $t_0 < t_1 < \dots$ such that $r < u(s_{t_k})$. By discarding finitely many elements of the sequence, we may assume that $-t_0 < r$. The definition of u now implies that the set $R(s_{t_k})$ is not empty for each $k \in \mathbb{N}$, and hence we can take an $r_k \in R(s_{t_k})$ such that $r < r_k$.

Suppose first there exists some $k \in \mathbb{N}$ for which $r_k \in R_*(s_{t_k})$. In that case, $x \in O(s_{t_k}) \subseteq \bigcap_{n \in \mathbb{N}} \{g_n > r_k\}$. It follows that $r < r_k < g_n(x)$ for each $n \in \mathbb{N}$ and consequently that $r \leq f(x)$.

Otherwise, for each $k \in \mathbb{N}$ choose an $n_k \in \mathbb{N}$ such that $r_k \in R_{n_k}(s_{t_k})$. We have $x \in O(s_{t_k}) \subseteq \{g_{n_k} > r_k\}$ and hence $r < r_k < g_{n_k}(x)$. It is therefore enough to show that the sequence n_0, n_1, \dots is unbounded: for then the numbers $g_{n_0}(x), g_{n_1}(x), \dots$ form a sequence converging to $f(x)$, and we are able to conclude that $r \leq f(x)$.

We argue that the sequence n_0, n_1, \dots is unbounded. Assume the contrary. By passing to a subsequence, we can then assume that $n_0 = n_1 = \dots$. Now the sequence r_0, r_1, \dots is bounded, because $r < r_k \leq \inf\{g_0(y) : y \in O(s_{t_k})\} \leq g_0(x)$, for each $k \in \mathbb{N}$. Since only finitely many points in the range of g_{n_0} fall in the interval $[r, g_0(x)]$, only finitely many of the sets $\{g_{n_0} > r_k\} : k \in \mathbb{N}\}$ are distinct. Thus, at least two of these sets are the same, say $\{g_{n_0} > r_0\} = \{g_{n_0} > r_1\}$. But s_{t_1} is a minimal sequence satisfying $O(s_{t_1}) \subseteq \{g_{n_0} > r_1\}$, while s_{t_0} is a proper initial segment of s_{t_1} satisfying $O(s_{t_0}) \subseteq \{g_{n_0} > r_0\}$, contradicting $r_1 \in R_{n_1}(s_{t_1})$. \dashv

We conclude this section with a list of some properties of the limsup functions that follow easily from the above characterization.

COROLLARY 2.2. *The sum, the minimum, and the maximum of two limsup functions is a limsup function.*

We say that a collection \mathcal{C} of real-valued functions on X is *closed under pointwise limits from above* if for each sequence $f_0 \geq f_1 \geq \dots$ of functions in \mathcal{C} converging pointwise to a function f , the function f is an element of \mathcal{C} .

COROLLARY 2.3. *The set of limsup functions is the smallest collection of functions that (a) contains all lower semicontinuous functions and (b) is closed under pointwise limits from above.*

COROLLARY 2.4. *A uniform limit of limsup functions is a limsup function.*

COROLLARY 2.5. *A function f is of Baire class 1 if and only if both f and $-f$ are limsup functions.*

§3. A game for limsup functions. In this section we turn to the analysis of the game $\Gamma(f)$. Let us begin with the following observation.

LEMMA 3.1. *Player II has a winning strategy in $\Gamma(f)$ precisely when f is a limsup function.*

The result is obvious: the rules of the game $\Gamma(f)$ are designed so that any function $u : T \rightarrow \mathbb{R}$ witnessing that f is a limsup function is a winning strategy for Player II, and vice versa.

Unlike the eraser game (see [8]), as we will show below, the game $\Gamma(f)$ need not be determined.

Recall that a set $B \subseteq X$ is called a *Bernstein set* if neither B nor $X \setminus B$ contains a non-empty perfect set. Under the axiom of choice, every uncountable Polish space contains a Bernstein set; moreover, every uncountable analytic set (in a Polish space) contains a non-empty perfect set.

THEOREM 3.2. *If X is uncountable and $B \subseteq X$ is a Bernstein set, then $\Gamma(1_B)$ is not determined.*

PROOF. Notice first that B is not a Borel set: for if it were, either B or $X \setminus B$ would contain a non-empty perfect set. And since B is not Borel, the function 1_B is not a limsup function, and hence Player II has no winning strategy in $\Gamma(1_B)$.

Suppose that Player I does.

Let $E = \{v \in 2^{\mathbb{N}} : v_n = 1 \text{ for infinitely many } n \in \mathbb{N}\}$, where we write $2 = \{0, 1\}$. Notice that E is not a Σ_2^0 subset of $2^{\mathbb{N}}$. For it were, we would be able to express $2^{\mathbb{N}}$ as a countable union of meagre sets, contradicting the Baire category theorem.

Now, consider the continuous function $g : 2^{\mathbb{N}} \rightarrow X$ induced by Player I's winning strategy. Then, $g(E) \subseteq X \setminus B$ and $g(2^{\mathbb{N}} \setminus E) \subseteq B$. Since $g(E)$ is an analytic subset of X , it is either countable, or it contains a perfect subset. But $X \setminus B$ contains no perfect subset. Thus $g(E)$ is countable, hence a Σ_2^0 subset of X . But then, $E = g^{-1}(g(E))$ is a Σ_2^0 subset of $2^{\mathbb{N}}$, yielding a contradiction. ⊥

We now turn to a sufficient condition for Player I to have a winning strategy. This sufficient condition will also turn out to be necessary under various assumptions.

Recall that a set is called a *Cantor set* if it is homeomorphic to the classical middle-thirds Cantor set.

THEOREM 3.3. *Let $f : X \rightarrow \mathbb{R}$ be arbitrary and suppose that there is a number $r \in \mathbb{R}$ and a Cantor set $C \subseteq X$ such that, in the subspace topology of C , the set $C \cap \{f \geq r\}$ is meagre and dense. Then, Player I has a winning strategy in $\Gamma(f)$.*

PROOF. Let $Y = C \cap \{f \geq r\}$, and let $\{S_0, S_1, \dots\}$ be a cover of Y by closed nowhere dense subsets of C . We presently construct a winning strategy for Player I.

Fix some sequence v_0, v_1, \dots of Player II's moves.

Let $y(0)$ be any point of the set $Y \setminus S_0$. Notice that the set $C \setminus S_0$ is not empty because S_0 is nowhere dense in C , and $Y \setminus S_0$ is not empty since Y is dense in C . Set $m_0 = 0$.

Player I starts with a move $x_0 = y(0)_0$. Take an $n \in \mathbb{N}$ and suppose that Player I's moves x_0, \dots, x_n have been defined, along with a point $y(n) \in Y$ and a number $m_n \in \mathbb{N}$, such that

$$(x_0, \dots, x_n) = (y(n)_0, \dots, y(n)_n). \tag{3.1}$$

To define the next move of Player I, x_{n+1} , we distinguish two cases:

Case 1: $v_n > r - 2^{-m_n}$ and $O(x_0, \dots, x_n) \cap S_{m_n} = \emptyset$. Let $y(n+1)$ be any point of the set $(O(x_0, \dots, x_n) \cap Y) \setminus S_{m_n+1}$. Notice that the set $(O(x_0, \dots, x_n) \cap C) \setminus S_{m_n+1}$ is not empty because S_{m_n+1} is nowhere dense in C , and $(O(x_0, \dots, x_n) \cap Y) \setminus S_{m_n+1}$ is not empty since Y is dense in C . Let $m_{n+1} = m_n + 1$, and define Player I's move as $x_{n+1} = y(n+1)_{n+1}$.

Case 2: otherwise. In this case we let $y(n+1) = y(n)$, $m_{n+1} = m_n$, and define Player I's move as $x_{n+1} = y(n+1)_{n+1}$.

Notice that in either case (3.1) holds for $n + 1$. This completes the definition of Player I's strategy.

The intuition behind this definition could be explained as follows: Player I starts by zooming in on the point $y(0)$ chosen to be in Y but not in S_0 . Player I awaits a stage n where Player II would make a move $v_n > r - 1$, and where the set S_0 would be "excluded." As soon as such a stage is reached, Player I switches to an element $y(1)$, chosen to be in Y but not in S_1 . He then zooms in on $y(1)$, awaiting a stage where Player I would make a move $v_n > r - 1/2$, and where S_1 would be "excluded." As soon as such a stage occurs, Player I switches to an element $y(2)$ chosen to be in Y but not in S_2 . And so on.

We argue that Player I's strategy is winning.

Suppose first that $\limsup v_n \leq r - 2^{-m}$ for some $m \in \mathbb{N}$. Then, Case 1 occurs at most finitely many times. Let N be the last stage when Case 1 occurs (or $N = 0$ if Case 1 never occurs). Then, the point x produced by Player I equals to $y(N)$. We thus have $\limsup v_n \leq r - 2^{-m} < r \leq f(y(N)) = f(x)$.

Suppose now that $\limsup v_n \geq r$. We argue that Case 1 occurs infinitely many times. Suppose to the contrary and let N be the last stage when Case 1 occurs (or 0 if Case 1 never occurs). Then, $m_N = m_{N+1} = \dots$ and $y(N) = y(N + 1) = \dots = x$. There are infinitely many $n > N$ with $v_n > r - 2^{-m_N}$, and for each such n the neighborhood $O(x_0, \dots, x_n)$ of x has a point in common with S_{m_N} . This implies that $x \in S_{m_N}$. This, however, contradicts the choice of $y(N)$. This establishes that Case 1 occurs infinitely many times.

Let x be the point constructed by Player I. We argue that $x \in C \setminus Y$. In view of (3.1), x is a limit of the sequence $y(0), y(1), \dots$. Since each $y(n)$ is an element of the closed set C , so is x . To see that x is not an element of Y , suppose to the contrary. Then, $x \in S_m$ for some $m \in \mathbb{N}$. Since Case 1 occurs infinitely often, the sequence m_0, m_1, \dots runs through all natural numbers, so we can choose $n \in \mathbb{N}$ to be the largest number such that $m_n = m$. This choice implies that Case 1 occurs at stage n , and hence $O(x_0, \dots, x_n)$ is disjoint from S_{m_n} , leading to a contradiction.

It follows that x is not an element of $\{f \geq r\}$. Thus $\limsup v_n \geq r > f(x)$, which completes the proof. \dashv

REMARK 3.4. For an arbitrary set $H \subseteq X$ the existence of a Cantor set $C \subseteq X$ such that $C \cap H$ is meager and dense in C is equivalent to the existence of a Cantor set $C \subseteq X$ such that $C \cap H$ is countable and dense in C . This either follows from Theorems 3.3 and 3.6 applied to 1_H and $r = \frac{1}{2}$, or can also be proved directly by a standard Cantor scheme construction.

If the range of the function f does not contain a strictly increasing sequence, the condition of Theorem 3.3 is both sufficient and necessary for Player I to have a winning strategy. The proof relies on a Kechris–Louveau–Woodin separation theorem [6, Theorem 21.22].

For a set $R \subseteq \mathbb{R}$, and function $f : X \rightarrow \mathbb{R}$ define the game $\Gamma_R(f)$ similarly to $\Gamma(f)$, but allowing Player II to choose v_i 's from R instead of the whole real line.

LEMMA 3.5. *Let $R \subseteq \mathbb{R}$ and $f : X \rightarrow R$. Then, Player I has a winning strategy in the game $\Gamma_R(f)$ if and only if Player I has a winning strategy in $\Gamma(f)$.*

PROOF. It is straightforward to check that if Player I has a winning strategy in $\Gamma(f)$, then the restriction of this strategy is winning for Player I in $\Gamma_R(f)$.

Conversely, fix a winning strategy σ_R of Player I for the game $\Gamma_R(f)$. For each $n \in \mathbb{N}$ define $F_n : \mathbb{R} \rightarrow R$ so that

$$\forall y \in \mathbb{R}, n \in \mathbb{N}: |F_n(y) - y| < d(y, R) + \frac{1}{n} \tag{3.2}$$

holds. Now define Player I’s strategy σ in $\Gamma(f)$ as follows: let $\sigma(\emptyset) = \sigma_R(\emptyset)$, and let

$$\sigma(x_0, v_0, x_1, v_1, \dots, x_n, v_n) = \sigma_R(x_0, F_0(v_0), x_1, F_1(v_1), \dots, x_n, F_n(v_n)) \tag{3.3}$$

whenever $n \in \mathbb{N}$ and $(x_0, \dots, x_n) \in T$.

It remains to check that σ is a winning strategy for Player I in $\Gamma(f)$. Fix a run $x_0, v_0, x_1, v_1, \dots$ of the game $\Gamma(f)$ consistent with σ , i.e., such that for each $n \in \mathbb{N}$, $\sigma(x_0, v_0, \dots, x_n, v_n) = x_{n+1}$. Then, (3.3) implies that for each n

$$\sigma_R(x_0, F_0(v_0), x_1, F_1(v_1), \dots, x_n, F_n(v_n)) = x_{n+1}, \tag{3.4}$$

and as σ_R is a winning strategy for Player I in $\Gamma_R(f)$, we obtain that

$$f(x_0, x_1, x_2, \dots) \neq \limsup_{n \rightarrow \infty} F_n(v_n).$$

We have to check that $f(x_0, x_1, x_2, \dots) \neq \limsup_{n \rightarrow \infty} v_n$. First, if $\limsup_{n \rightarrow \infty} v_n \notin R \supseteq \text{ran}(f)$, we are done. Otherwise $\limsup_{n \rightarrow \infty} v_n = r \in R$, therefore for each $\varepsilon > 0$, for all but finitely many k we have $v_k < r + \varepsilon$, thus (3.2) implies $F_k(v_k) < r + 2\varepsilon + \frac{1}{k}$ for these cofinitely many k ’s, therefore $\limsup_{k \rightarrow \infty} F_k(v_k) \leq r$. Since $r - \varepsilon < v_k$ holds for infinitely many k , this argument also shows that $r - 2\varepsilon - \frac{1}{k} < F_k(v_k)$ for infinitely many k too, thus

$$\limsup_{n \rightarrow \infty} v_n = r = \limsup_{n \rightarrow \infty} F_n(v_n) \neq f(x_0, x_1, x_2, \dots),$$

as desired. ◻

THEOREM 3.6. *Consider a function $f : X \rightarrow \mathbb{R}$ such that the range of f contains no infinite strictly increasing sequence. If Player I has a winning strategy in $\Gamma(f)$, then there is a number $r \in \mathbb{R}$ and a Cantor set $C \subseteq X$ such that the set $C \cap \{f \geq r\}$ is countable and dense in C .*

PROOF. Define R to be the closure of the range of f . Then, it is easy to verify that R contains no infinite strictly increasing sequence. Hence, the usual order $>$ of the reals is a well ordering of R . Let ρ be the order type of $(R, >)$, and let $\alpha \mapsto r_\alpha$ be the bijective map from ρ to R such that $r_\alpha > r_\beta$ whenever $\alpha < \beta$. Notice that ρ is a countable ordinal.

Assume that Player I has a winning strategy in $\Gamma(f)$. Then, by Lemma 3.5 Player I also has a winning strategy in $\Gamma_R(f)$. Let σ_R be such a strategy. Let $g : R^\mathbb{N} \rightarrow X$ be the continuous function induced by σ_R . Here R is given its discrete topology; since R is countable, $R^\mathbb{N}$ is a Polish space. For each $r \in R$ let $L_r = \{v \in R^\mathbb{N} : \limsup_{t \rightarrow \infty} v_t = r\}$, and let $A_r = g(L_r)$. The set A_r is analytic. Moreover,

$$\{f = r\} \cap A_r = \emptyset. \tag{3.5}$$

Suppose that the function f fails to satisfy the conclusion of the theorem, that is, there is no number $r \in \mathbb{R}$ and Cantor set $C \subseteq X$ such that $C \cap \{f \geq r\}$ is countable

and dense in C . We obtain a contradiction by showing that there exists a limsup function $e : X \rightarrow R$ such that Player I has a winning strategy in the game $\Gamma(e)$. More precisely, we show that σ_R is a winning strategy for Player I in $\Gamma_R(e)$, which suffices by Lemma 3.5.

Note that a function $e : X \rightarrow R$ is a limsup function if and only if $\{e \geq r\}$ is a Π_2^0 set for each $r \in R$, using that R is closed.

We define recursively a sequence $(G_\alpha : \alpha < \rho)$ of Π_2^0 subsets of X such that $\{f \geq r_\alpha\} \subseteq G_\alpha$. Let $\beta < \rho$ be an ordinal such that the sets $(G_\alpha : \alpha < \beta)$ have been defined. In particular, notice that

$$\{f > r_\beta\} \subseteq \bigcup_{\alpha < \beta} \bigcap_{\gamma: \alpha \leq \gamma < \beta} G_\gamma. \tag{3.6}$$

Since f fails to satisfy the condition of the theorem, there exists no Cantor set $C \subseteq X$ such that $C \cap \{f \geq r_\beta\}$ is countable and dense in C . This implies (using [6, Theorem 21.22]) that $\{f \geq r_\beta\}$ can be separated from any disjoint analytic subset of X by a Π_2^0 set. Consider the set

$$A_{r_\beta} \setminus \bigcup_{\alpha < \beta} \bigcap_{\gamma: \alpha \leq \gamma < \beta} G_\gamma. \tag{3.7}$$

It is analytic, since β is a countable ordinal. Moreover, it is disjoint from $\{f \geq r_\beta\}$ as can be seen from (3.5) and (3.6). Hence there exists a Π_2^0 subset G_β of X containing $\{f \geq r_\beta\}$ and disjoint from (3.7). Thus $\{f \geq r_\beta\} \subseteq G_\beta$ and

$$G_\beta \cap A_{r_\beta} \subseteq \bigcup_{\alpha < \beta} \bigcap_{\gamma: \alpha \leq \gamma < \beta} G_\gamma. \tag{3.8}$$

This concludes the recursive definition of the sequence $(G_\alpha : \alpha < \rho)$.

Now, for an arbitrary $\beta < \rho$ define the set

$$E_\beta = \bigcap_{\gamma: \beta \leq \gamma < \rho} G_\gamma.$$

This is a Π_2^0 set, since ρ is a countable ordinal. Moreover, $\{f \geq r_\beta\} \subseteq E_\beta$ for each $\beta < \rho$, hence $X = \bigcup_{\beta < \rho} E_\beta$. In view of (3.8), we have

$$E_\beta \cap A_{r_\beta} \subseteq \bigcup_{\alpha < \beta} E_\alpha. \tag{3.9}$$

Define $e : X \rightarrow R$ by letting $e(x) = r_\beta$, where $\beta < \rho$ is the least ordinal such that $x \in E_\beta$. Then, $\{e \geq r_\beta\} = E_\beta$ for each $\beta < \rho$, so e is a limsup function. Moreover, $\{e = r_\beta\}$ equals the set $E_\beta \setminus \bigcup_{\alpha < \beta} E_\alpha$, which is disjoint from A_{r_β} by (3.9). This shows that Player I's strategy σ_R remains winning in the game $\Gamma_R(e)$, yielding the desired contradiction. \dashv

Next we show that the above necessary and sufficient condition for the existence of a winning strategy for Player I also holds if we assume that f is sufficiently definable, e.g., Borel measurable.

We say that the function f is *semi-Borel* if for each $r \in \mathbb{R}$, the set $\{f \geq r\}$ is co-analytic.

THEOREM 3.7. *Let $f : X \rightarrow \mathbb{R}$ be semi-Borel. If Player I has a winning strategy in $\Gamma(f)$, then there is a number $r \in \mathbb{R}$ and a Cantor set $C \subseteq X$ such that the set $C \cap \{f \geq r\}$ is countable and dense in C .*

PROOF. If for each $r \in \mathbb{R}$, $\{f \geq r\}$ is a Π_2^0 set, then f is a limsup function, hence Player II has a winning strategy in $\Gamma(f)$, a contradiction. Hence $\{f \geq r\}$ is not a Π_2^0 set for some $r \in \mathbb{R}$. Then, the Hurewicz theorem (see, e.g., [6, Theorem 21.18]) implies that there is a Cantor set C such that the set $C \cap \{f \geq r\}$ is countable and dense in C . \dashv

COROLLARY 3.8. *If f is semi-Borel, then the game $\Gamma(f)$ is determined.*

PROOF. If for each $r \in \mathbb{R}$, $\{f \geq r\}$ is a Π_2^0 set, then f is a limsup function, hence Player II has a winning strategy in $\Gamma(f)$. Otherwise, $\{f \geq r\}$ is not a Π_2^0 set for some $r \in \mathbb{R}$, and as above, the Hurewicz theorem implies that there is a Cantor set C such that the set $C \cap \{f \geq r\}$ is countable and dense in C , therefore Player I has a winning strategy by Theorem 3.3. \dashv

Next we will show that in general the condition of Theorem 3.3 is not equivalent to the existence of a winning strategy for Player I. More precisely, we will show in Corollary 3.10 that the restriction on the range of f in Theorem 3.6 is optimal; if $R \subseteq \mathbb{R}$ contains an infinite strictly increasing sequence, then there exists a function $f : \mathbb{N}^{\mathbb{N}} \rightarrow R$ such that Player I has a winning strategy in $\Gamma(f)$, and $C \cap \{f \geq r\}$ is either uncountable or empty for each $r \in \mathbb{R}$ and each Cantor set $C \subseteq \mathbb{N}^{\mathbb{N}}$.

THEOREM 3.9. *There exists a function $f : \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}$ such that Player I has a winning strategy in $\Gamma(f)$, and $C \cap \{f \geq r\}$ is uncountable for each $r \in \mathbb{R}$ and each Cantor set $C \subseteq \mathbb{N}^{\mathbb{N}}$.*

PROOF. First note that every Cantor set can be written as a disjoint union of uncountably many (in fact, continuum many) Cantor sets, since it is well-known that a Cantor set is homeomorphic to 2^I for every countably infinite set I , in particular to $2^{\mathbb{N} \times \mathbb{N}}$, which is homeomorphic to $2^{\mathbb{N}} \times 2^{\mathbb{N}} = \bigcup_{c \in 2^{\mathbb{N}}} (\{c\} \times 2^{\mathbb{N}})$.

This implies that if H is an arbitrary set, then in order to show that $C \cap H$ is uncountable for each Cantor set $C \subseteq \mathbb{N}^{\mathbb{N}}$, it suffices to show that $C \cap H \neq \emptyset$ for each Cantor set $C \subseteq \mathbb{N}^{\mathbb{N}}$.

Let $X = \mathbb{N}^{\mathbb{N}}$ and let $\varphi : X \rightarrow \mathbb{N} \cup \{+\infty\}$ be given by $\varphi(x) = \limsup_{n \rightarrow \infty} x_n$. We first argue that there exists a function $f : X \rightarrow \mathbb{N}$ such that (a) $f(x) \neq \varphi(x)$ for each $x \in X$, and (b) $C \cap \{f \geq r\} \neq \emptyset$ for each $r \in \mathbb{R}$ and each Cantor set $C \subseteq X$. We will then show that condition (a) implies that Player I has a winning strategy in $\Gamma(f)$, while we already argued that (b) implies that $C \cap \{f \geq r\}$ is uncountable for each $r \in \mathbb{R}$ and each Cantor set $C \subseteq \mathbb{N}^{\mathbb{N}}$.

Let $(r_\alpha : \alpha < \mathfrak{c})$, $(z_\alpha : \alpha < \mathfrak{c})$, and $(C_\alpha : \alpha < \mathfrak{c})$ be enumerations of the real numbers, of the points of X , and of the Cantor subsets of X , respectively. We define the pairs $(z_\alpha, f(z_\alpha)) \in X \times \mathbb{N}$ recursively as follows. Take an ordinal $\alpha < \mathfrak{c}$ and suppose that $(z_\beta, f(z_\beta))$ has been defined for every $\beta < \alpha$. Let z_α be any point of $C_\alpha \setminus \{z_\beta : \beta < \alpha\}$. Define $f(z_\alpha)$ to be the smallest natural number such that $f(z_\alpha) \geq r_\alpha$ and $f(z_\alpha) \neq \varphi(z_\alpha)$. To complete the definition of f , for each point $x \in X \setminus \{z_\beta : \beta < \mathfrak{c}\}$ let $f(x)$ be the smallest natural number such that $f(x) \neq \varphi(x)$.

Now we show that Player I has a winning strategy in $\Gamma(f)$. Using Lemma 3.5, it is enough to show that Player I has a winning strategy in $\Gamma_{\mathbb{N}}(f)$. Let Player I start by playing $x_0 = 0$. To a move $v_n \in \mathbb{N}$ of Player II in round n , Player I responds with $x_{n+1} = v_n$. Then, for a run $x_0, v_0, x_1, v_1, \dots$ of the game, it holds that $\varphi(x) = \limsup_{n \rightarrow \infty} x_n = \limsup_{n \rightarrow \infty} v_n$. Since $f(x) \neq \varphi(x)$ holds, we have $f(x) \neq \limsup_{n \rightarrow \infty} v_n$, hence the run is won by Player I. \dashv

COROLLARY 3.10. *If $R \subseteq \mathbb{R}$ contains an infinite strictly increasing sequence, then there exists a function $f : \mathbb{N}^{\mathbb{N}} \rightarrow R$ such that Player I has a winning strategy in $\Gamma(f)$, and $C \cap \{f \geq r\}$ is either uncountable or empty for each $r \in \mathbb{R}$ and each Cantor set $C \subseteq \mathbb{N}^{\mathbb{N}}$.*

PROOF. Let $i : \mathbb{N} \rightarrow R$ be a strictly increasing map. Let $f_0 : \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{N}$ be a function as in Theorem 3.9, that is, such that Player I has a winning strategy in $\Gamma(f_0)$, and $C \cap \{f_0 \geq r\}$ is uncountable for each $r \in \mathbb{R}$ and each Cantor set $C \subseteq \mathbb{N}^{\mathbb{N}}$. We claim that the function defined as $f = i \circ f_0$ works. Clearly, $f : \mathbb{N}^{\mathbb{N}} \rightarrow R$, and it is also clear that $C \cap \{f \geq r\}$ is either uncountable or empty for each $r \in \mathbb{R}$ and each Cantor set $C \subseteq \mathbb{N}^{\mathbb{N}}$, hence we only have to show that Player I has a winning strategy in $\Gamma(f)$. By Lemma 3.5 it suffices to check that Player I has a winning strategy in $\Gamma_{i(\mathbb{N})}(f)$. Let σ_0 be a winning strategy for Player I in $\Gamma(f_0)$, and define for each $n \in \mathbb{N}$

$$\sigma_{i(\mathbb{N})}(x_0, v_0, \dots, x_n, v_n) = \sigma_0(x_0, i^{-1}(v_0), \dots, x_n, i^{-1}(v_n)). \tag{3.10}$$

Since i is order-preserving, it is easy to check that $\sigma_{i(\mathbb{N})}$ is a winning strategy for Player I in $\Gamma_{i(\mathbb{N})}(f)$. \dashv

Next we state another result of similar sort. We will strengthen the above counterexamples by showing that such an f can have a co-analytic graph, but on the other hand we have to sacrifice that the range is countable. Note that the complexity of the graph of f is optimal, since if the graph of a function is analytic, then it is well-known that the function is actually Borel measurable, hence by Theorem 3.7 it cannot be a counterexample, and similarly, the range cannot be countable, since it is easy to show that a function with co-analytic graph and countable range is semi-Borel.

Recall that the statement “ $V = L$ ” is the *Axiom of Constructibility* due to K. Gödel. It is known that it is consistent with *ZFC*, and that it implies the Continuum Hypothesis.

THEOREM 3.11. *Assume $V = L$. Then, there exists a function $f : \mathbb{N}^{\mathbb{N}} \rightarrow \mathbb{R}$ with co-analytic graph such that Player I has a winning strategy in $\Gamma(f)$, and $C \cap \{f \geq r\}$ is uncountable for each $r \in \mathbb{R}$ and each Cantor set $C \subseteq \mathbb{N}^{\mathbb{N}}$.*

PROOF. Let $X = \mathbb{N}^{\mathbb{N}}$. Let $q(0), q(1), \dots$ be an enumeration of the rational numbers, and let $\varphi : X \rightarrow \mathbb{R} \cup \{+\infty\}$ be given by $\varphi(x) = \limsup_{n \rightarrow \infty} q(x_n)$. In order to construct $f : X \rightarrow \mathbb{R}$ with co-analytic graph, we use a result of Vidnyánszky [13, Theorem 1.3]. Let $B_1 = \{(C, t) \in \mathcal{K}(X) \times \mathbb{R} : C \text{ is a Cantor set}\}$, where $\mathcal{K}(X)$ is the family of non-empty compact sets in X equipped with the Hausdorff metric. Let $B_2 = \mathbb{R}$ and $B = B_1 \sqcup B_2$ be the disjoint union of B_1 and B_2 making B a subset of the Polish space $(\mathcal{K}(X) \times \mathbb{R}) \sqcup \mathbb{R}$. Let $i : X \rightarrow \mathbb{R}$ be a Borel bijection, $M = \mathbb{R}^2$,

and let

$$F_1 = \left\{ (A, (C, r), (y, t)) \in M^{\leq \omega} \times B_1 \right. \\ \left. \times M : y \in i(C) \setminus \text{pr}_1(\text{ran}(A)), t \geq r, t \neq \varphi(i^{-1}(y)) \right\},$$

where $\text{pr}_1(\text{ran}(A))$ is the projection of the range of the sequence A onto the first coordinate. Let

$$F_2 = \left\{ (A, y', (y, t)) \in M^{\leq \omega} \times B_2 \times M : t \neq \varphi(i^{-1}(y)), y' \notin \text{pr}_1(\text{ran}(A)) \Rightarrow y' = y, \right. \\ \left. y' \in \text{pr}_1(\text{ran}(A)) \Rightarrow y \notin \text{pr}_1(\text{ran}(A)) \right\},$$

and let $F = F_1 \sqcup F_2 \subseteq M^{\leq \omega} \times B \times M$.

We now check that the conditions of Vidnyánszky’s theorem are satisfied. First, a non-empty compact set $C \subseteq \mathbb{N}^{\mathbb{N}}$ is a Cantor set if and only if it is perfect. Using [6, Exercise 4.31] one can easily see that B_1 is a Borel subset of $\mathcal{K}(X) \times \mathbb{R}$. Therefore B is a Borel subset of $(\mathcal{K}(X) \times \mathbb{R}) \sqcup \mathbb{R}$. The set F_1 is clearly co-analytic, and since $A \in M^{\leq \omega}$ is a countable sequence, conditions of the form $y' \in \text{pr}_1(\text{ran}(A))$ are Borel. Therefore F_2 is even Borel, making $F = F_1 \sqcup F_2$ co-analytic. For each $(A, b) \in M^{\leq \omega} \times B$, no matter whether $b \in \mathcal{K}(X) \times \mathbb{R}$ or $b \in \mathbb{R}$, the section

$$F_{(A,b)} = \{(y, t) \in M : (A, b, (y, t)) \in F\}$$

contains $\{x_1\} \times \{t : t \geq x_2\}$ for some $(x_1, x_2) \in \mathbb{R}^2$, hence it is cofinal in the Turing degrees (for this notion, see Definition 1.1 of [13]). Therefore the conditions of the theorem are satisfied.

The conclusion of the theorem assures that there is a co-analytic set $G \subseteq M = \mathbb{R}^2$ and enumerations $B = \{b_\alpha : \alpha < \omega_1\}$, $G = \{g_\alpha : \alpha < \omega_1\}$ and for every $\alpha < \omega_1$ a sequence $A_\alpha \in M^{\leq \omega}$ that is an enumeration of $\{g_\beta : \beta < \alpha\}$ such that $g_\alpha \in F_{(A_\alpha, b_\alpha)}$ for every $\alpha < \omega_1$. We note here that the assumption $V = L$ implies the continuum hypothesis.

First we check that G is the graph of a function with domain \mathbb{R} . Notice that for $\beta < \alpha$, if $g_\alpha = (y_1, t_1)$ and $g_\beta = (y_2, t_2)$, then $y_1 \neq y_2$. Indeed, $g_\alpha \in F_{(A_\alpha, b_\alpha)}$ implies that $y_1 \notin \text{pr}_1(\text{ran}(A_\alpha))$, and since $y_2 \in \text{pr}_1(\text{ran}(A_\alpha))$, $y_1 \neq y_2$ easily follows. To see that for each $y \in \mathbb{R}$, $(y, t) \in G$ for some $t \in \mathbb{R}$, let $\alpha < \omega_1$ be chosen with $b_\alpha = y \in B_2$. Then, either $y \in \text{pr}_1(\text{ran}(A_\alpha))$ and we are done, or g_α is chosen to be (y, t) for some $t \in \mathbb{R}$. Therefore G is indeed a graph of a function with domain \mathbb{R} .

Now we define the function $f : X \rightarrow \mathbb{R}$ in the following way: for each $(y, t) \in G$, let $f(i^{-1}(y)) = t$. Clearly, the graph of f is $(i, \text{id})^{-1}(G)$, hence it is co-analytic.

We now show that the defined function f has properties (a) $f(x) \neq \varphi(x)$ for each $x \in X$, and (b) $C \cap \{f \geq r\} \neq \emptyset$ for each $r \in \mathbb{R}$ and each Cantor set $C \subseteq X$. Then, we will show that (a) implies that Player I has a winning strategy in $\Gamma(f)$. The proof that (b) implies that $C \cap \{f \geq r\}$ is uncountable for each $r \in \mathbb{R}$ and each Cantor set C is exactly the same as in the proof of Theorem 3.9.

To show (a), let $(x, t) \in X \times \mathbb{R}$ be a pair with $(i(x), t) = g_\alpha \in G$. Then, $g_\alpha \in F_{(A_\alpha, b_\alpha)}$ implies $t \neq \varphi(x)$, hence $f(x) = t \neq \varphi(x)$. To show (b), let $C \subseteq X$ be a Cantor set and let $r \in \mathbb{R}$. Let $\alpha < \omega_1$ be the ordinal with $b_\alpha = (C, r)$. Then, for

$g_\alpha = (y, t)$, using again that $g_\alpha \in F_{(A_\alpha, b_\alpha)}$, $y \in i(C)$ and $t \geq r$, hence $i^{-1}(y) \in C$ and $f(i^{-1}(y)) = t \geq r$.

It remains to show that Player I has a winning strategy in $\Gamma(f)$. Let Player I start by playing $x_0 = 0$. To a move v_n of Player II, Player I responds with an $x_{n+1} \in \mathbb{N}$ chosen to be the smallest natural number satisfying $|v_n - q(x_{n+1})| \leq 2^{-n}$. Then, for a run $x_0, v_0, x_1, v_1, \dots$ of the game, it holds that $\varphi(x) = \limsup_{n \rightarrow \infty} v_n$. Since $f(x) \neq \varphi(x)$, the run is won by Player I. \dashv

We note that the assumption $V = L$ cannot be simply dropped from the above theorem. Indeed, it can be derived using the standard proof that Projective Determinacy implies that the Hurewicz theorem holds for all projective sets; moreover, if the graph of f is projective, then so is $\{f \geq r\}$ for every $r \in \mathbb{R}$. Thus one could derive an analogue to Theorem 3.7 under Projective Determinacy, assuming only that f has a projective graph.

Despite all the partial results above, we still do not know the answer to the following interesting question.

QUESTION 3.12. *For which $f : X \rightarrow \mathbb{R}$ does Player I have a winning strategy in $\Gamma(f)$?*

§4. A game for Baire class 1 functions. Recall the definition of the game $\Gamma'(f)$ from the Introduction. Corollary 2.5 immediately yields the following result:

COROLLARY 4.1. *Player II has a winning strategy in $\Gamma'(f)$ if and only if Player II has winning strategies in both games $\Gamma(f)$ and $\Gamma(-f)$, if and only if f is of Baire class 1.*

New we turn to the existence of a winning strategy for Player I.

Let $C \subseteq X$ be a closed set, and consider the restriction of f to C . The oscillation of $f|_C$ at a point $x \in C$ is defined as

$$\text{osc}_f(C, x) = \inf_{\substack{s \in T: \\ x \in O(s)}} \sup_{y, z \in O(s) \cap C} |f(y) - f(z)|.$$

LEMMA 4.2. *Suppose that there is a closed set $C \subseteq X$ such that the oscillation of $f|_C$ is bounded away from zero: $\inf_{x \in C} \text{osc}_f(C, x) > 0$. Then, Player I has a winning strategy in $\Gamma'(f)$.*

PROOF. Assume that $\text{osc}_f(C, x) \geq 5\varepsilon > 0$ for each $x \in C$. We will first describe a strategy of Player I and then we will show that it is a winning strategy. To define the moves of Player I in a particular run, we will use recursion to define natural numbers $n_0 < n_1 < n_2 < \dots$ and sequences $s_0, s_1, s_2, \dots \in T$ (these may depend on the moves of Player II).

Let $n_0 = 0$, and let s_0 be the empty sequence. Suppose that, for some even number $k \in \mathbb{N}$, Player I's moves prior to the stage n_k have been defined.

Let $s_k \in T$ denote the sequence of Player I's moves prior to the stage n_k . Define $\alpha_k = \sup\{f(x) : x \in O(s_k) \cap C\}$, and choose a point $x(k) \in O(s_k) \cap C$ so that $\alpha_k - \varepsilon < f(x(k))$. Starting with the stage n_k , Player I produces his moves using the point $x(k)$, that is, he plays $x_n = x(k)_n$ at a stage $n \geq n_k$. He continues doing so until the first stage, say $n_{k+1} > n_k$, that Player II makes a move $(v_{n_{k+1}}, w_{n_{k+1}})$ such

that $|v_{n_{k+1}} - f(x(k))| < \varepsilon$. If no such stage occurs, then Player I goes on using the point $x(k)$ to make his moves until the end of the game.

Let $s_{k+1} \in T$ denote the sequence of moves produced by Player I prior to the stage n_{k+1} . Define $\beta_{k+1} = \inf\{f(x) : x \in O(s_{k+1}) \cap C\}$, and choose a point $x(k+1) \in O(s_{k+1}) \cap C$ so that $f(x(k+1)) < \beta_{k+1} + \varepsilon$. Starting with the stage n_{k+1} , Player I produces the moves using $x(k+1)$, that is he plays $x_n = x(k+1)_n$ at a stage $n \geq n_{k+1}$. He continues doing so until the first stage, say $n_{k+2} > n_{k+1}$, that Player II makes a move $(v_{n_{k+2}}, w_{n_{k+2}})$ such that $|w_{n_{k+2}} - f(x(k+1))| < \varepsilon$. If no such stage occurs, then Player I goes on using the point $x(k+1)$ until the end of the game.

We show that the strategy thus defined is winning.

Suppose first that only finitely many stages n_0, n_1, \dots occur, the last one being n_k . For concreteness, suppose that k is even. In this case Player I uses the point $x(k)$ to generate his moves until the end of the game. Moreover, there is no $n > n_k$ such that $|v_n - f(x(k))| < \varepsilon$. This implies that $\limsup v_n \neq f(x(k))$, and hence the run is won by Player I. Likewise, if the last one of the sequence n_0, n_1, \dots is the stage n_{k+1} where k is even, then Player I generates the point $x(k+1)$, and there is no $n > n_{k+1}$ such that $|w_n - f(x(k+1))| < \varepsilon$. Therefore $\liminf w_n \neq f(x(k+1))$, and hence the run is won by Player I.

Suppose that infinitely many stages n_0, n_1, \dots occur. From the above definitions we get for each even $k \in \mathbb{N}$

$$\begin{aligned} v_{n_{k+1}} &> f(x(k)) - \varepsilon \\ &> \alpha_k - 2\varepsilon \\ &= (\alpha_k - \beta_{k+1}) + \beta_{k+1} - 2\varepsilon \\ &\geq (\alpha_k - \beta_{k+1}) + f(x(k+1)) - 3\varepsilon \\ &\geq (\alpha_k - \beta_{k+1}) + w_{n_{k+2}} - 4\varepsilon. \end{aligned}$$

Let $\alpha_{k+1} = \sup\{f(x) : x \in C \cap O(s_{k+1})\}$. Since the sequence s_{k+1} extends s_k , we have $\alpha_k \geq \alpha_{k+1}$. By the assumption, the oscillation of $f|_C$ at the point $x(k+1) \in C$ is at least 5ε , hence $\alpha_{k+1} - \beta_{k+1} \geq 5\varepsilon$. Combining these facts we obtain that for each even $k \in \mathbb{N}$ it holds that $v_{n_{k+1}} \geq w_{n_{k+2}} + \varepsilon$. This, however, means that $\limsup v_n > \liminf w_n$, implying a win for Player I. -1

REMARK 4.3. The above construction of the winning strategy for Player I is similar to that in [8]. In both cases Player I zooms in on a particular element of C until Player II triggers a switch to another element. The main difference is that here Player I undergoes two alternating types of switches: even switches are different from the odd ones. An odd switch, say $(k+1)$ st (where k is even) is triggered when Player II makes a move such that v_n is close to $f(x(k))$. Player I reacts by switching to a point $x(k+1)$ with a low value of f . Even switches, say $(k+2)$ nd, are triggered when Player II makes a move such that w_n is close to $f(x(k+1))$. Player I reacts by switching to a point $x(k+2)$ of C with a high value of f .

THEOREM 4.4. *Let $f : X \rightarrow \mathbb{R}$ be an arbitrary function. The game $\Gamma'(f)$ is determined.*

PROOF. If f is a function in Baire class 1, then Player II has a winning strategy by Lemma 4.1. Suppose that f is not a function in Baire class 1. Then (see [9, Theorem 2 and Remark 1, p. 395]) there exists a non-empty closed set $K \subseteq X$ such

that the set of discontinuity points of $f|_K$ contains an open subset of K . Using the Baire category theorem and the arguments as in [8, p. 9] one can show that there is a non-empty closed set $C \subseteq K$ such that the oscillation of $f|_C$ is bounded away from zero. The preceding lemma then implies that Player I has a winning strategy. \dashv

REMARK 4.5. It is not completely clear which results of the paper use the countability of A in an essential way. It seems to us that almost all results go through without the assumption that A is countable, and the only really problematic issues are the applications of the Hurewicz theorem and the Kechris–Louveau–Woodin theorem in the proofs of Theorems 3.6–3.8.

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