

# A GENERALIZATION OF DIFFERENCE SETS

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**1. Introduction.** A  $(v, k, \lambda)$  *difference set*  $D$  is a set of  $k$  distinct residues  $\{a_1, a_2, \dots, a_k\}$  modulo  $v$  such that every residue  $b \not\equiv 0 \pmod{v}$  can be expressed in exactly  $\lambda$  ways in the form  $b \equiv a_i - a_j \pmod{v}$ . With each difference set we may associate a binary periodic sequence  $(s_1, s_2, \dots)$  with  $s_i = 1$  if  $i \pmod{v}$  is in  $D$ , and  $s_i = 0$  otherwise. Since this sequence is periodic of period  $v$ , we need only consider one cycle from the sequence. Such cycles we agree to call (binary) *difference cycles*. Difference cycles (equivalently, difference sets) have been studied intensively (2, 4). They have important applications to digital communications, mainly because they have *2-level autocorrelation*. In this paper we shall point out certain other (equivalent) properties of difference cycles which seem susceptible to immediate generalization, but show that these generalizations are vacuous.

We wish to thank Professor E. S. Selmer for suggesting this problem.

**2. Motivation.** If  $s$  is a difference cycle, then the defining property of difference sets tells us that the number of ordered pairs  $(s_i, s_{i+b})$  from  $s$  (subscripts taken modulo  $v$ ) of the form  $(1, 1)$  is  $\lambda$  for all values of  $b \not\equiv 0 \pmod{v}$ . More generally, let the number of ordered pairs  $(s_i, s_{i+b})$  from  $s$  of the form  $(\epsilon_1, \epsilon_2)$  be denoted by  $p_{\epsilon_1, \epsilon_2}(b)$ . Thus  $p_{1,1}(b) = \lambda$  for all  $b \not\equiv 0 \pmod{v}$ . A simple enumeration now shows that, in addition,

$$p_{0,1}(b) = p_{1,0}(b) = k - \lambda, \quad p_{0,0}(b) = v - 2k + \lambda$$

whenever  $b \not\equiv 0 \pmod{v}$ . For let us represent  $s$  and its  $b$ th translate

$$s_b = (s_b, s_{b+1}, \dots, s_v, s_1, \dots, s_{b-1})$$

schematically as below:

$$\begin{array}{r}
 s: \quad \overbrace{1 \ 1 \ \dots \ 1 \ 1 \ 1 \ \dots \ 1}^k \ \overbrace{0 \ 0 \ \dots \ 0 \ 0 \ 0 \ \dots \ 0}^{v-k} \\
 s_b: \quad \overbrace{1 \ 1 \ \dots \ 1 \ 0 \ 0 \ \dots \ 0}^{\lambda} \ \overbrace{0 \ 0 \ \dots \ 0}^{k-\lambda} \ \overbrace{1 \ 1 \ \dots \ 1 \ 0 \ 0 \ \dots \ 0}^{k-\lambda} \ \overbrace{0 \ 0 \ \dots \ 0}^{v-2k+\lambda}
 \end{array}$$

Since exactly  $\lambda$  of the ones from  $s$  match up with ones from  $s_b$ , the remaining  $k - \lambda$  ones from  $s$  must be paired with zeros from  $s_b$ . This shows that  $p_{1,0}(b) = k - \lambda$ . The other relations may be verified similarly.

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Received November 12, 1965. This paper presents the results of one phase of research carried out at the Jet Propulsion Laboratory, California Institute of Technology, under Contract No. NAS 7-100, sponsored by the National Aeronautics and Space Administration.

Now let  $s$  be any binary cycle of length  $v$  (not necessarily associated with a difference set). The (unnormalized) autocorrelation of  $s$ ,  $R_s(b)$ , has been defined as follows **(2)**:  $R_s(b) = A_s(b) - D_s(b)$  where  $A_s(b)$  is the number of agreements between  $s$  and  $s_b$ , i.e. the number of components in which  $s$  and  $s_b$  have the same entry, and  $D_s(b)$  is the number of disagreements. With the notation introduced above,

$$R_s(b) = p_{1,1}(b) + p_{0,0}(b) - p_{1,0}(b) - p_{0,1}(b).$$

Our remarks then show in particular that for a difference cycle  $s$ ,  $R_s(b)$  is independent of  $b$  if  $b \not\equiv 0 \pmod{v}$ , and so  $s$  has two-level autocorrelation. (Of course  $R_s(0) = v$ .) Conversely, it is easy to show that any cycle with two-level autocorrelation is associated with a difference set **(2)**. In what follows, we shall refer to the trivial cycles of the form  $(1, 0, 0, \dots, 0)$  or  $(0, 1, 1, \dots, 1)$  or a translate of one of these as *pulses*. They correspond to the trivial difference sets with  $k = 1$  or  $k = v - 1$ .

**3. Generalization.** The preceding discussion motivates the following formal generalization of a difference cycle which depends neither on the notion of autocorrelation function nor on group theory. If  $s = (s_1, s_2, \dots, s_v)$  is any binary cycle, and if  $b$  is an integer satisfying  $0 < b < v$ , we define a *bigram*  $M(s, b)$  as follows:

$$M(s, b) = \{(s_i, s_{i+b}) : i = 1, 2, \dots, v\},$$

*multiplicity included.* We have seen that if  $s$  is a binary difference cycle, then  $M(s, b) = M(s, b')$  (equality means that the two collections contain the same pairs with the same multiplicity) whenever  $0 < b < v$ ,  $0 < b' < v$ . More generally, we define an *m-gram* for any *n-ary* cycle  $s = (s_1, s_2, \dots, s_v)$  as follows:

*Definition.* Let  $b = (b_1, b_2, \dots, b_{m-1})$  be an ordered  $(m - 1)$ -tuple of integers with  $0 < b_1 < b_2 < \dots < b_{m-1} < v$ . We define the *m-gram*  $M(s, b)$  to be the collection

$$\{(s_i, s_{i+b_1}, s_{i+b_2}, \dots, s_{i+b_{m-1}}), i = 1, 2, \dots, v\},$$

*multiplicity included.* If  $M(s, b) = M(s, b')$  for all such  $b$  and  $b'$ , we say that  $s$  is *m-tuply regular*.

Thus in the new terminology, an ordinary difference cycle becomes a doubly-regular binary cycle. We have formally generalized in two directions; we allow both the degree of regularity of  $s$  and the number of symbols in  $s$  to increase. We remark that *m-tuply regular unary* cycles and *singly regular n-ary* cycles are trivial, and that a pulse binary cycle is *m-tuply regular* in a trivial way. It is surprising that binary difference cycles and the above trivial examples are the only examples of *m-tuply regular n-ary* cycles possible. We prove this result now.

**4. Non-existence.**

**THEOREM 1.** *If  $s$  is an  $m$ -tuply regular  $n$ -ary cycle, then one of the three alternatives below holds:*

- (1)  $n = 2$  and either  $m = 2$  or  $s$  is a pulse,
- (2)  $m = 1$ ,
- (3)  $n = 1$ .

The *proof* is in two parts: (1) we assume that  $m > 2$  and conclude that  $n = 1$  unless  $s$  is a pulse; (2) we assume that  $n > 2$  and show that  $m = 1$ .

Suppose, then, that  $s$  is an  $m$ -tuply regular  $n$ -ary cycle with  $m > 2$ . It is clearly sufficient to prove this is impossible for  $n = 2$ , since if  $n > 2$  we may identify certain of the symbols to obtain an  $m$ -tuply regular binary cycle. We now digress in order to place the problem in a wider context.

*Definition (3):* Let  $S$  be a set of  $v$  distinct objects. A *tactical configuration*  $C = C[k, m, \lambda, v]$  is a collection of  $b$  subsets (called *blocks*)  $B_i, i = 1, 2, \dots, b$ , of  $S$  such that each block contains exactly  $k$  objects from  $S$ , and each (unordered)  $m$ -tuple from  $S$  occurs in exactly  $\lambda$  blocks.  $C$  is *symmetric* if  $b = v$ , and if each object in  $S$  occurs in exactly  $k$  blocks.

Our plan is to show that the existence of an  $m$ -tuply regular binary cycle implies the existence of a symmetric  $C[k, m, \lambda, v]$ . Theorem 2 shows that such configurations are trivial; we prove Theorem 2 first.

**THEOREM 2.** *There are no non-trivial symmetric  $C[k, m, \lambda, v]$  configurations, if  $m \geq 3$ . (Trivial means that  $k$  does not satisfy  $m \leq k \leq v - m$ .)*

*Proof.* It is clear that a  $C[k, m, \lambda, v]$  configuration is also a  $C[k, m', \lambda, v]$  configuration for  $m' \leq m$ , since each unordered  $m'$ -tuple from  $S$  is a subset of exactly  $\binom{v - m'}{m - m'}$  unordered  $m$ -tuples from  $S$ , and so each  $m'$ -tuple occurs

$$\lambda \binom{v - m'}{m - m'} / \binom{k - m'}{m - m'}$$

times in the configuration. Consequently it will be sufficient to prove Theorem 2 for  $m = 3$ .

For the moment let  $\lambda = \lambda_3$ , and let  $\lambda_2$  represent the number of times each unordered pair from  $S$  occurs in  $C$ . Then counting in two different ways the number of times a triple involving a given pair occurs in the configuration, we see that

$$(1) \quad \lambda_2(k - 2) = \lambda_3(v - 2).$$

Note that (1) holds for any  $C[k, 3, \lambda, v]$  configuration, symmetric or not.

If  $C$  is symmetric, let us count in two ways the number of times a pair involving a given element occurs:

$$(2) \quad k(k - 1) = \lambda_2(v - 1).$$

We now perform the standard trick (see **1**) of deleting from  $C$  an arbitrary block, and all objects occurring in that block. Since  $C$  is in particular a symmetric block design, the derived design  $C'$  will also be a  $C [k', 3, \lambda', v']$  configuration (but no longer symmetric) with  $k' = k - \lambda_2, \lambda' = \lambda'_3 = \lambda_3, \lambda'_2 = \lambda_2, v' = v - k$ . Equation (1) will now apply to the derived parameters; i.e.,

$$(3) \quad \lambda_2(k - \lambda_2 - 2) = \lambda_3(v - k - 2).$$

Combining (3) with (1), we see that  $\lambda_2^2/k = \lambda_3$ . Thus  $\lambda_2/k = (k - 2)(v - 2)$  and so, from (2),  $(k - 1)/(v - 1) = (k - 2)/(v - 2)$ , which implies that  $k = v$ . But  $k = v$  is a trivial design, and this completes the proof of Theorem 2.

To complete the first part of the proof of Theorem 1, it remains to show that the existence of an  $m$ -tuply regular binary cycle ( $m \geq 3$ ) which is not a pulse implies the existence of a non-trivial symmetric  $C [k, m, \lambda, v]$  configuration. Let  $k$  be the number of ones in the cycle. First of all, it is clear that if  $1 < k < m$ , then no such cycle exists, since for certain  $b$ 's,  $M(s, b)$  will contain

$$\overbrace{(1 \ 1 \ \dots \ 1)}^k \overbrace{(0 \ 0 \ \dots \ 0)}^{m-k} ,$$

while others will not. Similarly  $v - m > k > v - 1$  is impossible. Thus, except for pulses, all  $m$ -tuply regular binary sequences with  $k$  ones satisfy  $m \leq k \leq v - m$ .

If now  $s$  is an  $m$ -tuply regular binary cycle of length  $v$  (we assume the two symbols are 0 and 1), let  $S = \{a_1, a_2, \dots, a_v\}$  be any set containing  $v$  distinct objects. We define blocks  $B_i, i = 0, 1, 2, \dots, v - 1$ , as follows:  $a_j \in B_i$  if and only if  $s_{i+j} = 1$ . To show that these blocks form a symmetric  $C [k, m, \lambda, v]$  configuration, we need only verify the  $m$ -tuple condition.

Thus, let

$$(*) \quad (a_{i_1}, a_{i_2}, \dots, a_{i_m})$$

be an  $m$ -tuple from  $S$ , and assume that  $i_1 < i_2 < \dots < i_m$ . Let  $b_j = i_{j+1} - i_j, j = 1, 2, \dots, m - 1$ , and set  $\mathbf{b} = (b_1, b_2, \dots, b_{m-1})$ . Since  $s$  is  $m$ -tuply regular, the  $m$ -tuple  $(1, 1, \dots, 1)$  will occur in  $M(s, b)$  a certain number of times, say  $\lambda_m$ , and  $\lambda_m$  is independent of  $b$ . It is clear that if  $(s_{i_1}, s_{i_1+b_1}, \dots, s_{i_1+b_{m-1}})$  is such an  $m$ -tuple from  $M(s, b)$ , then  $B_{i_1}$  will contain the  $m$ -tuple  $(*)$  and conversely. Hence every  $m$ -tuple from  $S$  occurs in exactly  $\lambda_m$  blocks, and so the blocks  $B_i$  do form a (non-trivial) symmetric  $C [k, m, \lambda, v]$  configuration. But this is impossible by Theorem 2, and so every  $m$ -tuply regular binary cycle is a pulse. This completes the first part of the proof of Theorem 1.

Our attention has recently been drawn to the fact that the second half of Theorem 1 was proved independently by R. Titsworth (**5**) several years ago. (The reader who consults that report will see that Titsworth's "perfect" sequences are precisely the doubly regular sequences discussed here.) We

present here a new proof which makes use of the highly developed theory of difference sets.

In order to prove the second half of Theorem 1, we shall assume that  $s$  is a doubly regular  $n$ -ary cycle, and show that  $n > 2$  is impossible. It will be sufficient to prove that there are no doubly regular ternary cycles, since a doubly regular  $n$ -ary cycle ( $n > 3$ ) can be transformed into a doubly regular ternary cycle by a simple identification of certain symbols.

If  $s$  is a doubly regular ternary cycle in the symbols 0, 1, 2, let  $s$  contain  $k_0$  zeros,  $k_1$  ones, and  $k_2$  twos. We observe that each  $k_i$  must be  $\geq 2$ , since if (say)  $k_0 = 1$ , then some bigrams would contain (0, 1) but not (0, 2), while others would contain (0, 2) but not (0, 1). Let us now identify the symbols 0 and 1;  $s$  then becomes a doubly regular *binary* cycle (which is not a pulse by the observation above), and so the set  $D_2 = \{i: s_i = 2\}$  is a difference set. Similarly,  $D_0$  and  $D_1$  are also difference sets. But also  $D_{0,1} = \{i: s_i = 0 \text{ or } s_i = 1\}$ , being the complement of  $D_2$ , is a difference set. Also,  $D_{0,2}$  and  $D_{1,2}$ , defined similarly, are also difference sets.

We remark at this stage that  $D_0 \cup D_1 = D_{0,1}$ ,  $D_0 \cap D_1 = \emptyset$ . This means that if we could prove that the union of two disjoint difference sets is never a non-trivial difference set, the second part of Theorem 1 would follow as an immediate corollary. But although there is some evidence that this is so (see the discussion at the end of this paper), we are as yet unable to prove (or disprove) it. So we must use another route.

To continue with our proof, write

$$D_0 = \{a_1, a_2, \dots, a_{k_0}\} \quad \text{and} \quad D_1 = \{b_1, b_2, \dots, b_{k_1}\}.$$

We define the polynomials  $\theta_0$  and  $\theta_1$  as follows (cf. 4):

$$\begin{aligned} \theta_0(x) &\equiv x^{a_1} + x^{a_2} + \dots + x^{a_{k_0}} \pmod{x^v - 1}, \\ \theta_1(x) &\equiv x^{b_1} + x^{b_2} + \dots + x^{b_{k_1}} \pmod{x^v - 1}. \end{aligned}$$

Then since  $D_0$  and  $D_1$  are difference sets, we have, as in (4),

$$(4) \quad \begin{aligned} \theta_0(x)\theta_0(x^{-1}) &\equiv n_0 + \lambda_0 T(x) \pmod{x^v - 1}, \\ \theta_1(x)\theta_1(x^{-1}) &\equiv n_1 + \lambda_1 T(x) \pmod{x^v - 1}, \end{aligned}$$

where

$$\begin{aligned} n_0 &= k_0 - \lambda_0, & \lambda_0 &= k_0(k_0 - 1)/(v - 1), \\ n_1 &= k_1 - \lambda_1, & \lambda_1 &= k_1(k_1 - 1)/(v - 1), \end{aligned}$$

and

$$T(x) \equiv 1 + x + x^2 + \dots + x^{v-1} \pmod{x^v - 1}.$$

Now since  $s$  is doubly regular, the pair (1, 0) occurs in each bigram  $M(s, b)$  equally often, say  $\mu$  times. Hence every residue  $d \not\equiv 0 \pmod{v}$  can be written in exactly  $\mu$  ways in the form  $d \equiv a_i - b_j \pmod{v}$ . In terms of the  $\theta$ 's, this condition becomes

$$(5) \quad \theta_0(x)\theta_1(x^{-1}) \equiv -\mu + \mu T(x) \pmod{x^v - 1}.$$

If we multiply both sides of (5) by  $\theta_1(x)$ , and use (4), we obtain

$$(6) \quad n_1 \theta_0(x) + \mu \theta_1(x) \equiv (\mu k_1 + \lambda_1 k_0) T(x) \pmod{x^v - 1}.$$

(Observe that  $R(x)T(x) \equiv R(1)T(x) \pmod{x^v - 1}$ ; see (4).) But the left-hand side of (6) cannot contain powers of  $x$  higher than  $v - 1$ , and so

$$(7) \quad n_1 \theta_0(x) + \mu \theta_1(x) = (\mu k_1 - \lambda_1 k_0) T(x).$$

But this is impossible: the left-hand side of (7) cannot contain all the powers of  $x$  less than  $v$ , since some  $s_i$  are equal to 2. This contradiction completes the proof of Theorem 1.

**5. Conclusion.** We remark finally that there is a certain amount of arbitrariness in our definition of multiple regularity. With hindsight at least, we might regard the fact that the pairs (0, 1) and (1, 0) occur evenly distributed among the bigrams of a difference cycle as chance, peculiar to the case of binary cycles. We could then define multiple regularity by requiring only that  $m$ -tuples of the form  $(aa \dots a)$  be evenly distributed. If we had done this, the first part of the proof of Theorem 1 would still have worked, since we only needed the even distribution of  $(11 \dots 1)$  anyway. But a new proof of the second part of the theorem would be needed. In fact, it is easy to see that with the new definition the existence of a doubly-regular  $n$ -ary cycle is equivalent to the existence of  $n - 1$  disjoint difference sets  $(\text{mod } v)$  whose union is also a non-trivial difference set. If the parameters of these difference sets are  $(v, k_i, \lambda_i)$ ,  $i = 1, 2, \dots, n - 1$ , then an obvious necessary condition for their union to be a difference set is that  $v - 1 | k_i(k_i - 1)$  for  $i = 1, 2, \dots, n - 1$ , and  $v - 1 | k_n(k_n - 1)$ , where

$$k_n = v - \sum_{i=1}^{n-1} k_i.$$

More symmetrically, we require a solution to the number-theoretic problem of writing some  $v$  as  $v = k_1 + \dots + k_n$  so that  $v - 1 | k_i(k_i - 1)$  for  $i = 1, 2, \dots, n$ . For the case  $n = 3$ , computer search shows that there are exactly 51 solutions to this problem with  $v \leq 300$ ; and in every case it is quite easy to show that there is at least one  $i$  for which there is no  $(v, k_i, \lambda_i)$  difference set, so that the problem has no solution for  $n = 3$  and  $v \leq 300$ .

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