

## PLANAR MAPS ARE WELL LABELED TREES

ROBERT CORI AND BERNARD VAUQUELIN

In the theory of enumeration, the part devoted to the counting of planar maps gives rather surprising results. Of special interest to the combinatorialists is the conspicuous feature of counting numbers associated with families of maps as discussed in the papers of Tutte, Brown and Mullin. These formulas are by no means easy to prove; this is also a part of their charm. We are convinced therefore that these maps possess deep combinatorial properties. We discuss such properties in this paper.

The main result of this article is the construction of a bijection between planar maps and a special family of trees: their vertices are labeled with numbers that differ by at most 1 on adjacent vertices. These trees are said to be well labeled. Combinatorialists usually claim that theorems specifying the “geometry” of objects lead to enumerating formulas. Our result obeys this general rule: we can indeed obtain a new proof of the formula counting the number of rooted planar maps with  $m$  edges. Four different proofs of this formula are known to date. The first one was given by Tutte [14] using a bijection from maps with  $m$  edges onto maps with  $m$  vertices all of degree four and then using the counting formula for slicings [13]. A second proof by Tutte [15] consists in solving a functional equation verified by the generating formal power series of these numbers. A third one by Lehman [10] is established by encoding maps with parenthesis bracket systems. The fourth proof by Cori and Richard [3] consists in constructing a noncommutative generating formal power series for these numbers, then establishing an equation that it satisfies and finally building a solving process for this kind of equation.

Our paper, which we have tried to make self contained, is divided into five parts. In the first part we state the combinatorial definition of maps used by Edmonds [6], Jacques [8], Tutte and some others. This approach using a permutation enables us to define trees (usually called plane trees [4]) and well labeled trees. In the second part, we construct a bijection from planar maps onto well labeled trees in great detail and prove our main result (Theorem A).

In order to find Tutte’s formula we intend to encode well labeled trees with words that generalise parenthesis systems. Hence, the aim of the third part is to introduce essential notations and definitions on words, then to reconsider the classical construction of a tree by a parenthesis

---

Received May 24, 1979.

system using the representation of a tree by a permutation. The fourth part contains a study of the coding of well labeled trees. The reader who is familiar with formal language theory will notice that the set  $C$  of words encoding well labeled trees is a language which is the intersection of two context free languages but which is not itself context free. Finally, in the fifth section, we deduce the formula which enumerates the well labeled trees from a property of the language  $C$  (Theorem B) and hence we are able to enumerate planar maps.

## 1. Combinatorial maps.

1.1. *Definitions.* Let  $Z_m$  denote the set of non zero integers whose absolute value is at most  $m$ :

$$Z_m = \{-m, -m + 1, \dots, -2, -1, 1, 2, \dots, m\}.$$

For any  $\sigma$ , a permutation acting on  $Z_m$ ,  $z(\sigma)$  denotes the number of orbits of  $\sigma$  and  $\bar{\sigma}$  the permutation given by:

$$\bar{\sigma}(a) = \sigma(-a) \quad \text{for any } a \text{ in } Z_m.$$

A *combinatorial map* is a permutation  $\sigma$  such that the pair  $\{\sigma, \bar{\sigma}\}$  generates a transitive group on  $Z_m$ . This could be restated in the following manner: for any  $a$  and  $b$  in  $Z_m$ , there exist integers  $i_1, i_2, \dots, i_p$  such that:

$$\sigma^{i_1} \bar{\sigma}^{i_2} \dots \sigma^{i_p} \bar{\sigma}^{i_p}(a) = b.$$

The cycles of  $\sigma$  are called the *vertices* of the map, those of  $\bar{\sigma}$  its *faces*. A pair  $(a, -a)$  where  $a$  is in  $Z_m$  is an *edge*. The *genus*  $g(\sigma)$  of a map  $\sigma$  is given by:

$$g(\sigma) = \frac{1}{2}[m + 2 - z(\sigma) - z(\bar{\sigma})].$$

This number is known to be a non negative integer. A combinatorial map (or simply a *map*) is *planar* if its genus is zero, i.e.,  $z(\sigma) + z(\bar{\sigma}) = m + 2$  (Euler relation).

A planar map with only one face (i.e.,  $z(\bar{\sigma}) = 1$ ) is called a *tree*. A *labeled tree* is a tree in which a natural number (*label*) is assigned to each vertex. A tree will be said *well labeled* if labels which are assigned to adjacent vertices differ by at most one. More precisely: a mapping  $\epsilon$  from  $Z_m$  into  $\mathcal{N}$  is a well labeling of the tree  $\sigma$  if  $\epsilon(1) = 0$  and if, for any  $a$  in  $Z_m$ ,  $\epsilon(a) = \epsilon(\sigma(a))$  and  $|\epsilon(a) - \epsilon(-a)| \leq 1$ .

1.2. *Intuitive map representation.* Combinatorialists draw circles in order to represent the cycles of a permutation  $\sigma$ . The elements  $a, \sigma a, \dots, \sigma^k a$  of each cycle are set on the circle in a conventional way (clockwise for instance). Drawing an edge between  $a$  and  $-a$  for any  $a$  in  $Z_m$  will be our convention for representing the map  $\sigma$ . The reader familiar with the notion of topological map will find this map by con-

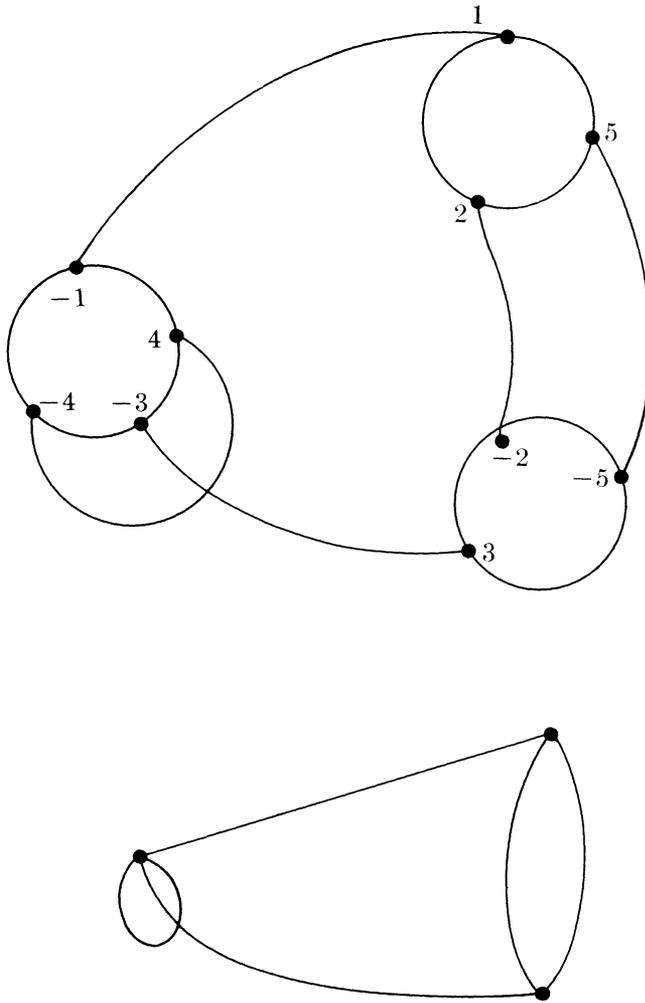


FIGURE 1. The representations of the combinatorial map:  $\sigma = (1, 5, 2)(3, -2, -5)(4, -3, -4, -1)$ .

tracting each circle to a point (Figure 1). Note that the underlying graph is always connected; this comes from the condition that  $\{\sigma, \bar{\sigma}\}$  generates a transitive group.

1.3. *Rooted maps.* In order to avoid some difficulties in the enumeration of maps coming from their symetries, rooted maps are usually introduced. Rooted maps are the classes of the following isomorphism relation:

Two maps  $\sigma$  and  $\sigma'$  are isomorphic if there exists a permutation  $\varphi$

such that:

$$\varphi(1) = 1, \varphi\sigma\varphi^{-1} = \sigma \quad \text{and, for any } a \text{ in } Z_m, \quad -\varphi(a) = \varphi(-a).$$

In a rooted map the *root vertex* and the *root edge* are the vertex and the edge which contains the element 1 of  $Z_m$ .

With this definition one finds 2 and 9 rooted planar maps with respectively 1 and 2 edges. The same number is found for well labeled trees (see Figure 2).

**2. The main theorem.** This section is devoted to the proof of the main theorem. We will often use two basic operations on permutations: the *closure* of a subset and the *restriction*. With these two operations we define what we call the *splitting*  $\hat{\sigma}$  of a map  $\sigma$ . We prove that  $\hat{\sigma}$  is a tree when  $\sigma$  is a planar map. In well labeling  $\hat{\sigma}$  we obtain then a one-to-one mapping from rooted planar maps onto rooted well labeled trees. In order to prove this fact we establish first a lemma: one can reconstruct a permutation  $\sigma$  on a set  $A$  from its action on a subset  $A'$  and its restriction to the complement  $A''$  of  $A'$ .

2.1. *General notations on permutations.* Let  $\sigma$  be a permutation acting on a set  $A$  and  $B$  be a subset of  $A$ . In the sequel  $\sigma^*B$  will denote the minimal subset of  $A$  containing  $B$  and closed by  $\sigma$ . Then the *closure* is defined by:

$$\sigma^*B = \{a \in A; \exists n \in \mathbf{N}, \exists b \in B \text{ such that } a = \sigma^n(b)\}.$$

If  $\sigma^*B = B$ ,  $B$  will be said to be *saturated* by  $\sigma$ . Lastly, the restriction  $\sigma|_B$  of  $\sigma$  to  $B$  is the permutation acting on  $B$  verifying:

$$\sigma|_B(b) \text{ is the first of } \sigma(b), \sigma^2(b) \dots \sigma^n(b) \text{ belonging to } B.$$

*Example.* Let:

$$\begin{aligned} \sigma = (1, -4, 4, 2, 5)(-2, -6, -7, 9)(8, -8, -5, -1) \\ \times (-3, 3, 7, -9)(6) \end{aligned}$$

be a permutation acting on  $Z_9$ . Let  $B = \{-2, -6, 6, -7, -9\}$ ; then  $\sigma^*B = \{-2, -6, -7, 9, 6, -9, -3, 3, 7\}$  and  $\sigma|_B = (-2, -6, -7)(6)(-9)$ .

2.2. *The splitting of a map.* For any map  $\sigma$  acting on  $Z_m$  we construct a sequence of subsets  $B_0, B_1, \dots, B_i, \dots, B_p$ , in the following manner:

$$\begin{aligned} B_0 &= \sigma^*\{1\} \quad \text{and for any } i \geq 0 \\ B_{2i+1} &= \sigma^*B_{2i} \setminus B_{2i} = \{b \in \sigma^*B_{2i}, b \notin B_{2i}\} \\ B_{2i+2} &= \sigma^*B_{2i+1} \setminus B_{2i+1}. \end{aligned}$$

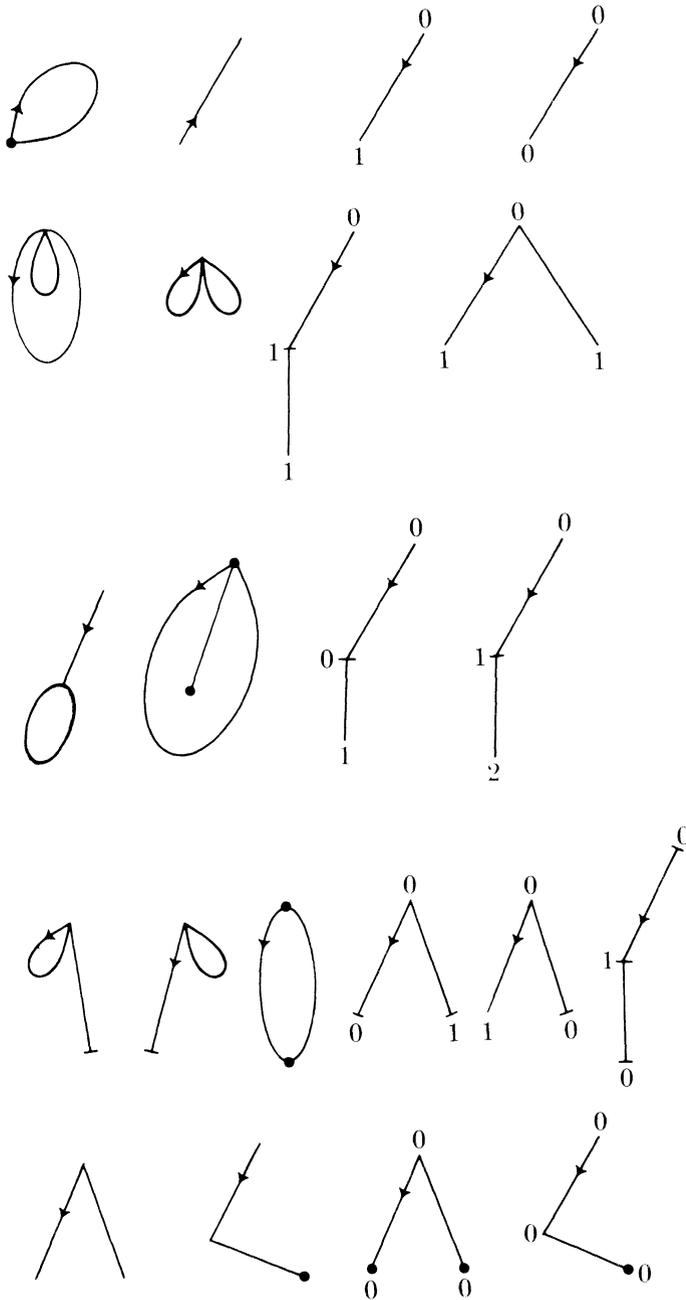


FIGURE 2. Rooted planar maps and well labeled trees with 1 and 2 edges.

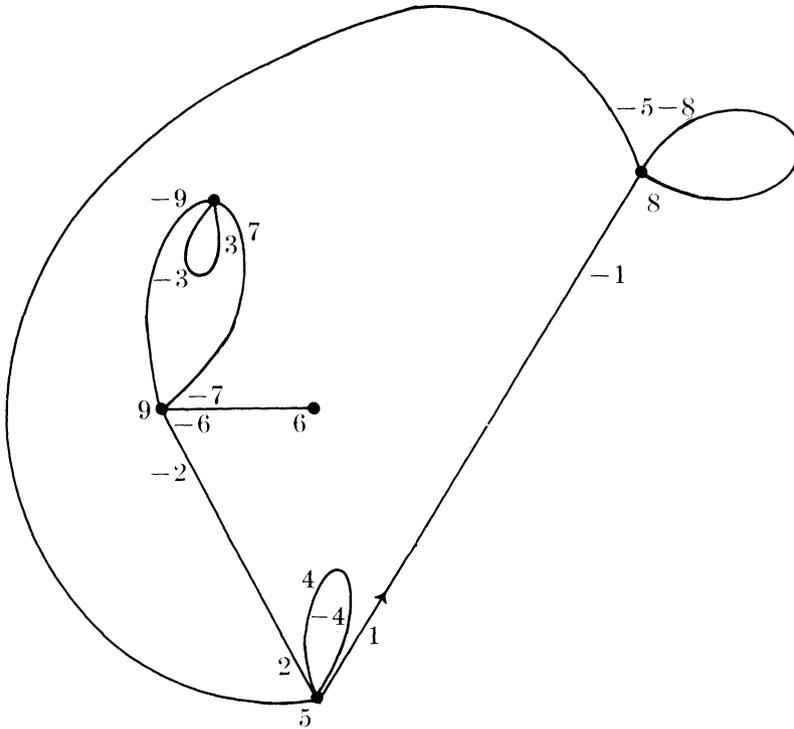


FIGURE 3. The map  $\sigma = (1, -4, 4, 2, 5) (-2, -6, -7, 9) (8, -8, -5, -1) (-3, 3, 7, -9) (6)$ .

*Example.* For the permutation  $\sigma$  given above, representing the map of figure 3:

$$\bar{\sigma} = (1, 8, -5)(-1, -4, 2, -6, 6, -7, -9, -2, 5)(3) \times (-3, 7, 9)(4)(-8)$$

$$B_0 = \{1, 8, -5\} \quad B_1 = \{-1, 4, -4, 5, 2, -8\}$$

$$B_2 = \{-9, -2, -7, 6, -6\} \quad B_3 = \{7, 9, 3, -3\}$$

$$B_4 = \emptyset = B_i \quad (i > 4).$$

PROPOSITION 2.1. *The permutation  $\sigma$  saturates  $B_{2i} \cup B_{2i+1}$ . Also  $B_{2i+1} \cup B_{2i+2}$  is saturated by  $\bar{\sigma}$ .*

*Proof.* From the definition of the  $B_i$  one has:

$$\sigma^*B_{2i} = B_{2i} \cup B_{2i+1} \quad \text{and} \quad \bar{\sigma}^*B_{2i+1} = B_{2i+1} \cup B_{2i+2}.$$

The result then follows from the fact that  $\sigma^*(\sigma^*(A)) = \sigma^*(A)$ .

PROPOSITION 2.2. *The  $B_j$ 's define a partition of  $Z_m$ .*

*Proof.* Let us first show that the  $B_j$ 's are mutually disjoint, by induction on  $j$ . As  $B_0$  and  $B_1$  are disjoint let us assume that  $B_0, B_1, \dots, B_j$  are disjoint and suppose that  $j$  is even (the case  $j$  odd is obtained by exchanging  $\sigma$  and  $\bar{\sigma}$ ).

Let  $b$  be an element of  $B_{2i+1}$ , then  $b = \sigma^k(a)$  where  $a$  belongs to  $B_{2i}$ . If  $b$  is in  $B_{2i'+1}$  (or in  $B_{2i'}$ ) with  $0 \leq i' < i$  as  $a = \sigma^{2m-k}(b)$ , then  $a$  would be in  $\sigma^*B_{2i'+1}$  (or in  $\sigma^*B_{2i'}$ )  $\subseteq B_{2i'} \cup B_{2i'+1}$  from Proposition 2.1, which contradicts the fact

$$(B_{2i'} \cup B_{2i'+1}) \cap B_{2i} = \emptyset.$$

Hence  $b$  does not belong to any  $B_p$  ( $p \leq 2i - 1$ ) no more it does to  $B_{2i}$  from the very definition of  $B_{2i+1}$ . The union of the subsets  $B_i$  is exactly  $Z_m$ : because of the transitivity of the group generated by  $\{\sigma, \bar{\sigma}\}$ , as for each  $b$  in  $Z_m$  there exist  $i_1, i_2, \dots, i_p$  such that if

$$b = \sigma^{i_1}\bar{\sigma}^{i_2} \dots \sigma^{i_p-1}\bar{\sigma}^{i_p}(1),$$

then  $b$  belongs to some  $B_n$  with  $n \leq p$ .

*Definition.* The splitting of a map  $\sigma$  is the permutation  $\hat{\sigma}$  given by

$$\begin{aligned} \hat{\sigma}(a) &= \sigma_{/B_{2i}}(a) \text{ if } a \text{ belongs to } B_{2i} \\ &= \bar{\sigma}_{/B_{2i+1}}(a) \text{ if } a \text{ belongs to } B_{2i+1}. \end{aligned}$$

For the above example  $\sigma$  we thus obtain:

$$\begin{aligned} \hat{\sigma} &= (1)(8, -5)(-1, -4, 2, 5)(4)(-8)(-6, -7, -2)(6)(-9) \\ &\quad \times (3)(-3, 7, 9). \end{aligned}$$

### 2.3. Property of the splitting operation.

**THEOREM 2.3.** *The splitting of a map is a map.*

We have to prove that the group generated by  $\{\hat{\sigma}, \bar{\hat{\sigma}}\}$  acts transitively on  $Z_m$ . For this, we are going to construct for each  $a$  and  $b$ , such that  $a = \sigma(b)$  or  $a = \bar{\sigma}(b)$ , a word  $f$  on the alphabet  $\{\hat{\sigma}, \bar{\hat{\sigma}}\}$  such that  $a = f(b)$ ; the result then follows by the transitivity of the group generated by  $\{\sigma, \bar{\sigma}\}$ .

1. Let us consider, first, the case where  $a$  belongs to the set  $B_n$  for which  $B_{n+1} = \emptyset$ , and let us suppose that  $n$  is even ( $n = 2i$ ), the case  $n$  odd being treated in the same manner, exchanging  $\sigma$  and  $\bar{\sigma}$ . If  $a = \sigma(b)$ ,  $B_{2i}$  is saturated by  $\sigma$ ; as  $\hat{\sigma}$  is equal to  $\sigma_{/B_{2i}}$  in  $B_{2i}$  then  $a = \hat{\sigma}(b)$  and thus  $f = \hat{\sigma}$ . If  $a = \bar{\sigma}(b)$ , then  $b = -\sigma^{-1}(a)$ . Now as above,  $\sigma^{-1}(a) = \hat{\sigma}^{-1}(a)$  hence  $b = -\hat{\sigma}^{-1}(a)$  and  $a = \hat{\sigma}(-b)$ .

2. Let us suppose that the result is true if  $a$  belongs to  $B_j$ , for all  $j > 1$ . Let  $a = \sigma b$  be in  $B_i$ , and let  $i = 2k$ . One can deal in the same way with the case  $a \in B_{2k+1}$ . If  $b$  is in  $B_{2k}$ , then  $a = \hat{\sigma}b$  and the result is proved. So,

let  $b$  be in  $B_{2k+1}$  and let us compute the sequence:

$$b_1 = b, b_2 = \sigma^{-1}b_1, \dots, b_{l+1} = \sigma^{-1}b_l;$$

all the elements of this sequence are in  $B_{2k} \cup B_{2k+1}$ ; the first one is in  $B_{2k+1}$  and the last one in  $B_{2k}$ ; then there exists  $p$  such that:

$$b_1, b_2, \dots, b_{p-1} \in B_{2k+1}, \quad b_p \in B_{2k}.$$

Then we have  $\hat{\sigma}(b_p) = a$  (as  $\hat{\sigma} = \sigma|_{B_{2k}}$  for the elements of  $B_{2k}$ ). By induction there exists  $f_l$  such that  $b_{l+1} = f_l(b)$  (where  $1 \leq l \leq p - 1$ ). Thus

$$a = \hat{\sigma}f_{p-1} \dots f_1(b).$$

If  $a = \bar{\sigma}(b)$  we have  $a = \sigma(-b)$  with  $a$  in  $B_{2k}$ ; since we have proved in the first part that there exists  $f$  such that  $a = f(-b)$ , the result is true with  $f$ .

**2.4. Splitting a planar map.** For any  $a$  in  $Z_m$ , let  $\epsilon(a)$  be equal to  $j$  if  $a$  belongs to  $B_j$ ; then  $\epsilon$  defines a labelling of  $\sigma$ ; moreover we have:

**PROPOSITION 2.4.** *The splitting of a planar map is a tree, well labeled by  $\epsilon$ .*

*Proof.* (a) Let us first determine the number of cycles of  $\hat{\sigma}$ . From its definition each  $B_j$  is saturated by  $\hat{\sigma}$ ; further:

$$\begin{aligned} z(\hat{\sigma}|_{B_{2i}}) &= z(\sigma|_{B_{2i}}) = z(\sigma|_{B_{2i} \cup B_{2i+1}}) \\ z(\hat{\sigma}|_{B_{2i+1}}) &= z(\bar{\sigma}|_{B_{2i+1}}) = z(\bar{\sigma}|_{B_{2i+1} \cup B_{2i+2}}). \end{aligned}$$

Taking the sum of these relations we obtain:

$$z(\sigma) = \sum_{i=0}^{[n/2]} z(\sigma|_{B_{2i} \cup B_{2i+1}}) + \sum_{i=1}^{[n/2]} z(\bar{\sigma}|_{B_{2i-1} \cup B_{2i}}).$$

Then,  $z(\hat{\sigma}) = z(\sigma) + z(\bar{\sigma}) - 1$ , as the second sum begins with  $i = 1$  and  $B_0$  is a cycle of  $\bar{\sigma}$ .

(b) The genus  $g(\hat{\sigma})$  is now given by:

$$\begin{aligned} g(\hat{\sigma}) &= \frac{1}{2}(2 + m - z(\hat{\sigma}) - z(\hat{\sigma})) \\ g(\hat{\sigma}) &= g(\sigma) + \frac{1}{2}[1 - z(\hat{\sigma})]. \end{aligned}$$

The map  $\sigma$  being planar,  $g(\sigma) = 0$ ; moreover the genus of a map is a natural number [7], thus  $z(\hat{\sigma})$  is necessarily equal to 1 and  $\hat{\sigma}$  is a tree, thus it is a planar map.

(c) To prove that this tree is well labeled by  $\epsilon$  we must verify that if  $b$  belongs to  $B_j$ , then  $-b$  is in one of the subsets  $B_{j-1}, B_j, B_{j+1}$ .

As  $B_{2i} \cup B_{2i-1}$  is saturated by  $\bar{\sigma}$  one has:

$$\bar{\sigma}B_{2i} \subset B_{2i-1} \cup B_{2i};$$

then, applying  $\sigma^{-1}$  to these sets and using the fact that  $B_{2k} \cup B_{2k+1}$  is saturated by  $\sigma$  (then also by  $\sigma^{-1}$ ) we obtain:

$$(1) \quad -B_{2i} \subset B_{2i-2} \cup B_{2i-1} \cup B_{2i} \cup B_{2i+1}.$$

Starting with  $\sigma B_{2i} \subset B_{2i} \cup B_{2i+1}$  and applying  $\bar{\sigma}^{-1}$  we obtain:

$$(2) \quad -B_{2i} \subset B_{2i-1} \cup B_{2i-1} \cup B_{2i} \cup B_{2i+2}.$$

The result for  $j$  even now follows by comparing (1) and (2). The same kind of proof can be given for  $j$  odd.

2.5. *A useful lemma.* We have associated to each rooted planar map a rooted well labeled tree having the same number of edges. We have now to show that this correspondance is one-to-one. For this we need the following lemma which allows us to “extend” a permutation.

*Definition.* Let  $\varphi$  be a mapping from  $A$  to  $A$ .  $\varphi$  admits a cycle if there exist  $a_1, a_2, \dots, a_p$  such that  $\varphi(a_i) = a_{i+1}$  ( $1 \leq i < p$ ) and  $\varphi(a_p) = a_1$ .

LEMMA. Let  $A'$  and  $A''$  be two disjoint subsets the union of which is  $A$ . Let  $\varphi$  be a one-to-one mapping from  $A'$  into  $A$  which admits no cycle in  $A'$ . Let  $\alpha$  be a permutation acting on  $A''$ . Then there exists a unique permutation

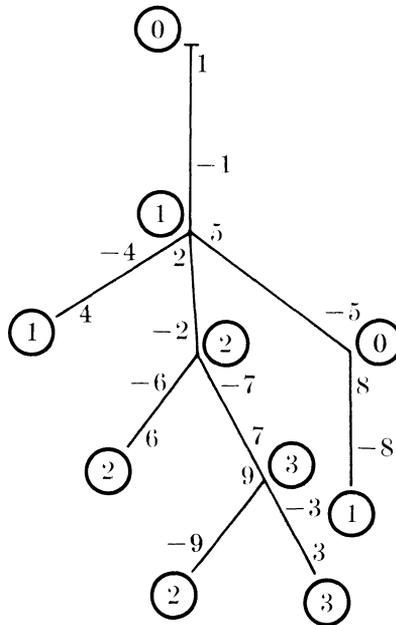


FIGURE 4. The well labeled tree  $\hat{\sigma}$  associated to the map  $\sigma$  of figure 3.

$\theta$  acting on  $A$  such that:

- (i)  $\theta|_{A''} = \alpha$ ,
- (ii)  $\theta(a) = \varphi(a)$  for every  $a$  in  $A'$ .

In the sequel we will denote by  $T_{A'',A'}(\alpha, \varphi)$  the permutation  $\theta$  which may be considered as an extension of  $\alpha$  by  $\varphi$ .

*Proof of the Lemma.* (by induction on the number of elements of  $A'$ )

If  $A'$  is empty then  $\theta = \alpha$  is the unique permutation.

If  $A'$  is not empty let  $b$  be an element of  $A'$  the image of which by  $\varphi$  is not in  $A'$  (such an element exists as  $\varphi$  admits no cycle). By induction, we obtain a unique permutation  $\theta'$  on  $A \setminus \{b\}$  such that  $\theta'|_{A''} = \alpha$  and  $\theta'(a) = \varphi(a)$  for  $a \in A', a \neq b$ . Let  $x$  be the image of  $b$  by  $\varphi$ .  $y = \theta'^{-1}(x)$  is necessarily an element of  $A''$  as  $\varphi$  is one-to-one. The permutation defined by:

$$\theta(b) = x, \theta(y) = b, \theta(a) = \theta'(a) \quad \text{for } a \neq b, a \neq y$$

is clearly the unique permutation which satisfies the conditions of the lemma.

2.6. THEOREM A. *The mapping  $\sigma \rightarrow (\hat{\sigma}, \epsilon)$  is a one-to-one correspondence between rooted planer maps and rooted well labeled trees.*

*Proof.* We have to reconstruct  $\sigma$  from  $(\hat{\sigma}, \epsilon)$ ; for this, we use the following remarks.

- (a) The subsets  $B_0, B_1, \dots, B_n$  are determined by:

$$a \in B_j \Leftrightarrow \epsilon(a) = j.$$

- (b) On  $B_n$ ,  $\hat{\sigma}$  is equal to  $\sigma$  or  $\bar{\sigma}$  according to the parity of  $n$ . For instance, if  $n$  is even then  $B_n$  is saturated by  $\sigma$  and  $\sigma(a) = \hat{\sigma}(a)$  for every  $a$  in  $B_n$ .

- (c) For any  $a$ ,  $\sigma^{-1}(a) = \bar{\sigma}^{-1}(a)$ .

- (d) For any  $i$ ,  $\sigma^{-1}|_{B_{2i}} = \hat{\sigma}^{-1}|_{B_{2i}}$  and  $\bar{\sigma}^{-1}|_{B_{2i-1}} = \hat{\sigma}^{-1}|_{B_{2i-1}}$ .

Let us now describe the reconstruction process if  $n$  is even (replace  $\sigma$  by  $\bar{\sigma}$  if  $n$  is odd).

$\sigma$  is first determined in  $B_n$  by Remark 2(b).

Remarks (c) and (d) imply that:

$$\bar{\sigma}^{-1}|_{B_{n-1} \cup B_n} = T_{B_{n-1}, B_n}(\hat{\sigma}^{-1}|_{B_{n-1}}, -\sigma^{-1}),$$

then the lemma of Section 2.5 allows the construction of  $\bar{\sigma}^{-1}$  on  $B_{n-1} \cup B_n$ .

Then remarks (c) and (d) again imply that:

$$\sigma^{-1}|_{B_{n-2} \cup B_{n-1}} = T_{B_{n-2}, B_{n-1}}(\hat{\sigma}^{-1}|_{B_{n-2}}, -\bar{\sigma}^{-1}),$$

which allows the construction of  $\sigma^{-1}$  on  $B_{2n-2} \cup B_{2n-1}$ , and so on. We

can thus claim that  $\sigma$  and  $\bar{\sigma}$  are uniquely determined by:

$$\begin{aligned} \hat{\sigma}(a) &= \sigma(a) \quad \text{if } a \in B_n \quad \text{and} \\ \bar{\sigma}^{-1}/_{B_{2i-1} \cup B_{2i}} &= T_{B_{2i-1} \cup B_{2i}}(\hat{\sigma}^{-1}/_{B_{2i-1}}, -\sigma^{-1}), \\ \sigma^{-1}/_{B_{2i} \cup B_{2i+1}} &= T_{B_{2i} \cup B_{2i+1}}(\hat{\sigma}^{-1}/_{B_{2i}}, -\bar{\sigma}^{-1}). \end{aligned}$$

**3. Coding trees by parenthesis systems.** In this part we begin by stating our notations concerning words written over an alphabet. We shall use this notion for the classical construction of the code of a rooted plane tree [9], [4]. We will briefly state the coding algorithm using the combinatorial definition of a tree.

3.1. *Parenthesis systems.* Let  $X$  be a finite set called *alphabet*; a *word* is a mapping from  $[n]$  (where  $[n] = \{1, 2, \dots, n\}$ ) into  $X$ . We denote by  $|f|$  the *length*  $n$  of  $f$ . One generally writes a word  $f$  by  $f = f(1)f(2) \dots f(n)$ ; for example if  $X = \{a, b\}$  and  $f(2i + 2) = a, f(2i + 1) = b$  for  $i = 0, 1$  then  $f$  is written  $f = baba$ . The *product* (or concatenation) of two words  $f$  and  $g$  of respective lengths  $n$  and  $m$  is the word  $h$  of length  $n + m$  given by:  $h(i) = f(i)$  if  $i \leq n$ , and  $h(i) = g(i - n)$  if  $n + 1 \leq i \leq n + m$ . If  $f = aba, g = ba$  is written  $ababa$ . The set  $X^*$  of all the words written on the alphabet  $X$  has a structure of free monoid the neutral element of which is the empty word 1. A word  $g$  is a *prefix* of the word  $f$  if one can find  $h$  such that  $f = gh$  (where  $h$  is a word, possibly empty).

Let us consider the alphabet  $X = \{x, \bar{x}\}$  and let  $\delta$  be the morphism from  $X^*$  into  $Z$  (the set of integers) given by:

$$\delta(x) = 1, \delta(\bar{x}) = -1 \quad \text{and} \quad \delta(f) = \sum_{i=1}^n \delta(f(i)).$$

*Definition.* A word  $f$  is a *parenthesis system* if:

- (i)  $\delta(f) = 0$
- (ii)  $\delta(g) \geq 0$  for every prefix  $g$  of  $f$ .

Let us denote by  $D$  the set of all parenthesis systems.

3.2. *Counting parenthesis systems.* Let  $D\bar{x}$  be the set of words obtained from those of  $D$  by right product of each one with  $\bar{x}$ . The following proposition [11] enables the counting of the words of length  $2n$  of  $D$ .

PROPOSITION 3.1. *For each word  $f$  in  $X^*$  such that  $\delta(f) = -p$  ( $p > 0$ ) there exist exactly  $p$  pairs  $(g_i, h_i)$  such that  $f = g_i h_i, h_i \neq 1$  and  $h_i g_i$  is a product of  $p$  elements of  $D\bar{x}$ .*

*Proof.* For every  $j$  ( $1 \leq j \leq q$ ) let  $f_j$  be the shortest prefix of the word  $f$  such that  $\delta(f_j) = -j$ ; then  $f$  may be written:  $f = u_1 u_2 \dots u_q v$  with  $f_i = u_i \dots u_i$ . The pairs  $g_i, h_i$  are then given by

$$g_i = u_1 \dots u_{q-p+1} \quad h_i = u_{q-p+i+1} \dots u_q v.$$

PROPOSITION 3.2. *The number of  $p$ -uples of parenthesis systems whose lengths sum to  $2n$  is:*

$$\frac{p}{2n + p} \frac{(2n + p)!}{n!(n + p)!}.$$

*Proof.* From Proposition 3.1, the number  $a_{n,p}$  of words of  $(D\bar{x})^p$  of length  $2n + p$  is equal to  $p/(2n + p)$  times the number of words  $f$  in  $X^*$  such that  $\delta(f) = -p$ . The fact that this last set contains all the words with  $n$  “ $x$ ” and  $(n + p)$  “ $\bar{x}$ ” implies

$$a_{n,p} = \frac{p}{2n + p} \frac{(2n + p)!}{n!(n + p)!}.$$

The correspondance

$$(f_1, f_2, \dots, f_p) \rightarrow (f_1\bar{x}f_2\bar{x}f_3 \dots f_p\bar{x})$$

is one-to-one, thus the proof of property follows from the fact that

$$|f_1\bar{x}f_2\bar{x} \dots f_p\bar{x}| = \sum_{i=1}^p |f_i| + p.$$

3.3. *Coding rooted trees.* Usually this coding is presented in the following intuitive manner: “Walk around the tree starting from the root and turning counter clockwise. Write  $x$  when going along an edge for the first time,  $\bar{x}$  when going along it again” (see figure 5). If a tree is defined as a combinatorial map  $\sigma$  such that  $z(\sigma) = m + 1$ ,  $z(\bar{\sigma}) = 1$  then the above algorithm should be made precise: first a canonical tree  $\sigma_0$  (isomorphic to  $\sigma$ ) is constructed from  $\sigma$ ; then the coding is obtained from  $\sigma_0$ .

Canonical tree: a tree  $\sigma$  is *canonical* if the permutation  $\bar{\sigma}$  satisfies:

- (1) If  $i, j$  are such that  $1 \leq i < j \leq n$  and  $\bar{\sigma}^i(1) > 0$ ,  $\bar{\sigma}^j(1) > 0$ , then  $\bar{\sigma}^i(1) > \bar{\sigma}^j(1)$ .
- (2) If  $i, j$  are such that  $\bar{\sigma}^i 1 = -\bar{\sigma}^j 1 > 0$  then  $i < j$ .

PROPOSITION 3.3. *For every rooted tree  $\sigma$  there exists a unique canonical tree  $\sigma_0$  isomorphic to  $\sigma$ .*

This isomorphism is obtained by renaming the elements of  $Z_m$  using  $\bar{\sigma}$ . For example, let us consider the tree  $\hat{\sigma}$  constructed in Section 2:

$$\hat{\sigma} = (1)(8, -5)(-1, -4, 2, 5)(4)(-8)(-6, -7, -2)(6)(-9) \\ (3)(-3, 7, 9).$$

One finds:

$$\hat{\sigma} = (1, -4, 4, 2, -6, 6, -7, 9, -9, -3, 3, 7, -2, 5, 8, -8, \\ -5, -1).$$

Hence,

$$\hat{\sigma}_0 = (1, 2, -2, 3, 4, -4, 5, 6, -6, 7, -7, -5, -3, 8, 9, -9, -8, -1)$$

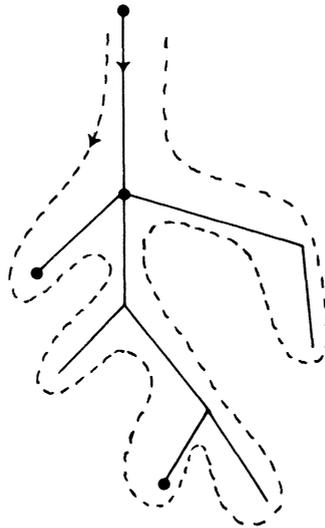


FIGURE 5. Coding:  $x\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}\bar{x}$  of the tree (1), (8, -5) (-1, -4, 2, 5)(4) (-8)(-6, -7, -2)(6)(-9)(3)(-3, 7, 9).

and

$$\hat{\sigma}_0 = (1)(2, 3, 8, -1)(-2)(-3, 4, 5)(-4)(-5, 6, 7)(-6)(-7) \times (-8, 9)(-9).$$

*Coding a canonical tree.* Let  $\sigma$  be a canonical tree. The coding  $\gamma(\sigma)$  of  $\sigma$  is the word  $f$  defined by:

$$f(i) = x \text{ if } \bar{\sigma}^{i-1}(1) \text{ is positive}$$

$$f(i) = \bar{x} \text{ if } \bar{\sigma}^{i-1}(1) \text{ is negative.}$$

Clearly condition (2) on a canonical tree implies that  $f$  belongs to  $D$ . Conversely given a word  $f$  in  $D$  we associate to it a tree  $\sigma$  (such that  $\gamma(\sigma) = f$ ) using the two following algorithms:

A1. For every prefix  $f$  of a word belonging to  $D$  we associate the subset  $A(f)$  of  $[n]$  (where  $n$  is the length of  $f$ ) by the recursive procedure:

$$A(1) = \emptyset$$

$$A(jx) = A(f) \cup \{|f|_x\}$$

(where  $|g|_x$  denotes the number of “ $x$ ” occurring in  $g$ )

$$A(j\bar{x}) = A(f) \setminus \text{Max } A(f).$$

If  $f = x\bar{x}\bar{x}\bar{x}\bar{x}$  then  $A(f) = \{2, 3, 5\}$ .

A2. For every  $f$  in  $D$ ,  $\gamma^{-1}(f)$  is the permutation  $\bar{\theta}$  acting on  $Z_m$  (where  $2m = |f|$ ) such that  $\theta$  is the circular permutation defined by:

$$\begin{aligned} \theta^i(1) &= \text{Max } A(f_{i+1}) & \text{if } f(i+1) = x \\ \theta^i(1) &= -\text{Max } A(f_i) & \text{if } f(i+1) = \bar{x} \end{aligned}$$

(where  $f_i$  denotes the prefix of length  $i$  of  $f$ ).

**4. The coding of well labeled trees.**

4.1. *Code description.* In the last section we defined the way to code a rooted tree; in the case of well labeled trees some information has to be added in the code word in order to find the labels again. This could be done by giving subscripts to the letters  $x$  and  $\bar{x}$  appearing in the code word. More precisely, walking around the infinite face of the tree (as for rooted trees) one of the letters  $x_1, x_2, x_3, \bar{x}_1, \bar{x}_2$  or  $\bar{x}_3$  is written according to the following rules:

—  $x_i$  (resp.  $x_2, x_3$ ) is written if the edge is traversed for the first time and if the initial vertex label is greater than (resp. equal to, less than) the terminal vertex one.

—  $\bar{x}$  ( $i = 1, 2$  or  $3$ ) is written if the edge is traversed for the second time and when  $x_i$  was written in the first encounter.

See figure for example.

Formally, the coding of a well labeled rooted tree is given by the following:

To each canonical well labeled rooted tree  $(\sigma, \epsilon)$  is associated the word  $f = \tilde{\gamma}(\sigma, \epsilon)$  on the alphabet  $X_3 = \{x_1, x_2, x_3, \bar{x}_1, \bar{x}_2, \bar{x}_3\}$  in the following way: let  $b_i = \sigma^{i-1}1$ ; then

$$\begin{aligned} f(i) &= x_{2+\epsilon(b_i)-\epsilon(-b_i)} & \text{if } b_i > 0 \\ f(i) &= \bar{x}_{2+\epsilon(-b_i)-\epsilon(b_i)} & \text{if } b_i < 0. \end{aligned}$$

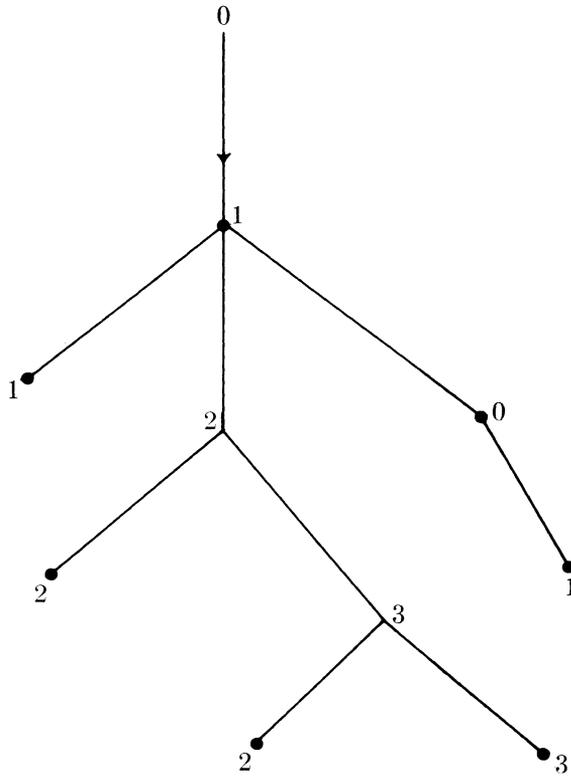
4.2. *The  $D_3$  language.* A first glance at the codes for well labeled rooted trees shows that they contain an equal number of  $x_i$  and  $\bar{x}_i$ ; moreover each  $x_i$  has a matching  $\bar{x}_i$ . This property comes from the fact that when walking around the tree and traversing the edge  $a_1$  and further the edge  $a_2$ , then  $a_2$  will be traversed again before  $a_1$  will be.

More formally, we have the following definition and proposition.

*Definition.* Let  $\rightarrow$  be the relation defined on the words of  $X_3^*$  by:

$$f \rightarrow g \Leftrightarrow f = f_1 x_i \bar{x}_i f_2 \quad \text{and} \quad g = f_1 f_2.$$

Let  $\xrightarrow{*}$  be the transitive closure of  $\rightarrow$ . We will say that  $f$  reduces to  $g$  if  $f \xrightarrow{*} g$ .



$x_1x_2\bar{x}_2x_1x_2\bar{x}_2x_1x_3\bar{x}_3x_2\bar{x}_2\bar{x}_1\bar{x}_1x_3x_1\bar{x}_1\bar{x}_3\bar{x}_1$ .

FIGURE 6. The code word of the labeled tree given in figure 4.

PROPOSITION 4.1. *If  $f$  is a code of a well labeled rooted tree, then  $f$  reduces to the empty word  $\|$ .*

Let  $D_3$  denote the set of words of  $X_3^*$  reducing to the empty word, let  $C$  be the set of codes of well labeled rooted trees. Then  $C$  is strictly included in  $D_3$  as  $x_3\bar{x}_3$  belongs to  $D_3$  but not to  $C$ .

The characteristic properties of the words in  $C$  will be given later. Let us establish first some properties of  $D_3$ .

*Definition.* A word  $f$  in  $D_3$  is prime if it has no prefix in  $D_3$  (other than  $\|$  or  $f$ ).

PROPOSITION 4.2. *Every word of  $D_3$  different from  $\|$  has a unique factorisation as a product of prime words.*

*Proof.* If  $f$  is not prime  $f = f_1g$  where  $f_1$  is prime; it is clear that  $g$  reduces to  $\|$  and repeating this decomposition to  $g$  one finds  $f = f_1f_2 \dots f_p$ .

PROPOSITION 4.3. *If  $f$  is prime then  $f = x_i g \bar{x}_i$  and  $g$  belongs to  $D_3$ .*

*Proof.* As  $f$  belongs to  $D_3$ , its first letter is necessarily an  $x_i$ ; this  $x_i$  has a matching  $\bar{x}_i$  and so  $f = x_i f_1 \bar{x}_i f_2$  where  $f_1$  and  $f_2$  reduce to  $\parallel$ . Clearly  $f_2 = \parallel$  otherwise  $x_i f_1 \bar{x}_i$  would be a left factor of  $f$  in  $D_3$ .

PROPOSITION 4.4. *The number of  $p$ -tuples of words in  $D_3$  whose lengths sum to  $2n$  is given by*

$$3^n \frac{p}{2n + p} \frac{(2n + p)!}{n!(n + p)!}.$$

*Proof.* The elements of  $D_3^p$  are obtained from those of  $D^p$  by choosing subscripts (1, 2, or 3) for each  $x$ , those of the  $\bar{x}$  being deduced from the  $x_i$  matching them. Then one gets the formula enumerating  $D_3^p$  multiplying by  $3^n$  the number of elements of  $D^p$  given by Proposition 3.2.

4.3. *Characterisation of the words of  $C$ .* Let  $\Delta$  be the mapping of  $X_3$  in  $\mathbf{Z}$  given by  $\Delta(x_1) = \Delta(\bar{x}_3) = 1, \Delta(x_2) = \Delta(\bar{x}_2) = 0, \Delta(\bar{x}_1) = \Delta(x_3) = -1$ . This mapping can be extended to a morphism of  $X_3^*$  into  $\mathbf{Z}$  (considered as a monoid for addition), denoted also by  $\Delta$  in the following way:

$$\begin{aligned} -\Delta(\parallel) &= 0 \\ -\Delta(fy) &= \Delta(f) + \Delta(y) \quad y \in X_3. \end{aligned}$$

PROPOSITION 4.5. *A word  $f$  belongs to  $C$  if and only if:*

- (1)  $f \in D_3$
- (2)  $\Delta(f') \geq 0$  for every prefix  $f'$  of  $f$ .

*Proof.* The code  $\tilde{\gamma}(\sigma, \epsilon)$  is obtained from  $\gamma(\sigma) = g$  in giving subscripts to the letters appearing in  $g$ . This way to give subscripts is in relation with the labeling  $\epsilon$  by:

$$\Delta(f') = \epsilon(\tilde{\sigma}^{f'} 1) \text{ for every prefix } f' \text{ of } f = \tilde{\gamma}(\sigma, \epsilon).$$

This relation enables the construction of  $\epsilon$  from the code  $f$  using  $\Delta(f')$ . The property  $\Delta(f') \geq 0$  ensures that labels are positive or null.

### 5. Counting well labeled trees.

5.1. *Main results.* The coding of well labeled trees enables us to find now the formula counting them. More precisely, we first state a combinatorial result (Theorem B) on the words of  $D_3$  which do not belong to  $C$ ; we then deduce as a corollary the enumerating formula. In order to prove Theorem B we establish four lemmas.

THEOREM B. *There exists a bijection  $\beta$  of  $D_3 \setminus C$  onto  $D_3 \times D_3 \times D_3$  satisfying:*

$$\text{if } \beta(f) = (f_1, f_2, f_3) \text{ then } |f| = |f_1| + |f_2| + |f_3| + 2.$$

COROLLARY. *The number  $a_m$  of well labeled rooted trees is given by:*

$$a_m = \frac{2 \cdot 3^m (2m)!}{m!(m + 2)!}.$$

*This is also the number of rooted planar map with  $m$  edges.*

*Proof* (of the corollary).  $a_m$  is also the number of words of  $C$  with length  $2m$ . From Theorem *B* this is the difference between the number of words of  $D_3$  with length  $2m$  and the number of triples of words in  $D_3$  whose lengths sum to  $2m - 2$ . These two numbers have been determined in Section 4 (Proposition 4.4). They are respectively equal to:

$$\frac{3^m (2m)!}{m!(m + 1)!} \quad \text{and} \quad \frac{3^m (2m + 1)!}{2m + 1 (m - 1)!(m + 2)!}.$$

The first part of the corollary is then obtained by the computation of their difference. The second part is a direct consequence of Theorem *A* (Section 2).

In order to prove Theorem *B* we have to associate three words of  $D_3$  to each word  $f$  of  $D_3 \setminus C$ . Lemma 1 will give us a decomposition of  $f$  as a product of three words. This decomposition is not a bijection onto  $D_3 \times D_3 \times D_3$ , but we may construct (Lemma 2) a bijection  $\beta_1$  of  $D_3 \setminus C$  onto  $D_3 \times D_3 \times D_3 \times \{1, 2, 3\}$ . Lemma 3 and Lemma 4 will then give combinatorial results on  $\mathbb{N}^3$  which will allow to prove the theorem.

#### 4.2. The mapping $\beta_1$ .

*Definition.* Let  $\partial$  be the mapping from  $X^*$  into  $Z$  given by:

$$\partial(f) = \text{Min} (\Delta f' \mid f' \text{ is a prefix of } f).$$

As the empty word 1 is always a prefix of  $f$  and as  $\Delta \parallel = 0$  we have that  $\partial f$  is a negative or zero number; moreover  $\partial f$  equals zero if and only if  $f$  belongs to  $C$ .

LEMMA 5.1. *Every word  $f$  in  $D_3 \setminus C$  has a unique factorisation as  $f = uvw$  where  $u, v, w$  are three words in  $D_3$  satisfying:*

- (1)  $v$  is a prime word
- (2)  $\partial u > \partial v$
- (3)  $\partial w \geq \partial v$ .

*Proof.* From Proposition 4.2,  $f$  can be decomposed as  $f = f_1 f_2 \dots f_p$ , where each  $f_i$  is a prime word. Let  $v = f_i$  be the leftmost with minimal image by  $\partial$ : we have thus either  $\partial f_j > \partial f_i$  or  $\partial f_j = \partial f_i$  and  $j \geq i$  in this case.

Then, denoting  $u = f_1 f_2 \dots f_{i-1}$  and  $w = f_{i+1} \dots f_p$  we obtain the decomposition of the lemma. The fact that  $\partial u > \partial v$  follows from  $\partial u f_i = \partial f_i$  as  $\Delta u = 0$  and as  $\partial f_i$  is minimal and strictly negative. The unicity is clear.

Let  $f$  be a word in  $D_3 \setminus C$ . From Lemma 5.1,  $f$  can be written as  $uvw$  and as  $v$  is a prime word  $v = x_i v' \bar{x}_i$  ( $v' \in D_3$ ). Let us denote  $\beta_1(f) = (u, v', w, i)$ . Lemma 5.1 could be stated in the following way:

LEMMA 5.2.  $\beta_1$  is a bijection from  $D_3 \setminus C$  onto  $\tilde{U}$ , the subset of elements  $(u, v, w, i)$  of  $D_3 \times D_3 \times D_3 \times \{1, 2, 3\}$  satisfying  $\partial u > \partial(x_i v \bar{x}_i)$  and  $\partial w \geq \partial(x_i v \bar{x}_i)$ .

5.3. A combinatorial property of  $\mathbb{N}^3$ .

LEMMA 5.3. Let  $\mathbb{N}$  be the set of natural numbers  $E_1, E_2, E_3$  the subsets of  $\mathbb{N}^3$  defined by:

$$E_1 = \{(a, b, c) \mid c + 1 < a\} \quad E_2 = \{(a, b, c) \mid a \leq b\} \\ E_3 = \{(a, b, c) \mid b \leq c\}.$$

Then the following subsets  $A = E_1 \cap \bar{E}_2$ ,  $B = E_2 \cap \bar{E}_3$ ,  $C = E_3 \cap \bar{E}_1$  form a partition of  $\mathbb{N}^3$  (where for any  $E, \bar{E} = \{x \in \mathbb{N}^3 \mid x \notin E\}$ ).

Proof. The fact that  $A, B, C$  are mutually disjoint could easily be verified. Now  $E_1 \cap E_2 \cap E_3 = \emptyset$  as  $(a, b, c) \in E_1 \cap E_2 \cap E_3$  implies  $c < a \leq b \leq c$ . Similarly  $\bar{E}_1 \cap \bar{E}_2 \cap \bar{E}_3 = \emptyset$ . The classical relation

$$\mathbb{N}^3 = (\bar{E}_1 \cap \bar{E}_2 \cap \bar{E}_3) \cup (E_1 \cap E_2 \cap E_3) \cup (E_1 \cap \bar{E}_2) \\ \cup (\bar{E}_1 \cap E_3) \cup (E_2 \cap \bar{E}_3)$$

finally gives the result.

LEMMA 5.4. Let  $U$  be the subset of  $\mathbb{N}^4$  given by  $U = U_1 \cup U_2 \cup U_3$  where:

$$U_1 = \{(a, b, c, d) \mid d = 1, a < b - 1, c \leq b - 1\} \\ U_2 = \{(a, b, c, d) \mid d = 2, a < b, c \leq b\} \\ U_3 = \{(a, b, c, d) \mid d = 3, a < b + 1, c \leq b + 1\}.$$

Then the mapping  $\varphi$  defined below is one-to-one from  $U$  onto  $\mathbb{N}^3$ :

$$\varphi(a, b, c, 1) = (b, c, a); \quad \varphi(a, b, c, 2) = (c, b, a) \\ \varphi(a, b, c, 3) = (c, a, b).$$

Proof. One verifies first that  $\varphi$  restricted to each of  $U_1, U_2, U_3$  is one-to-one. The relations  $\varphi(U_1) = A, \varphi(U_2) = B, \varphi(U_3) = C$  and Lemma 5.3 give the result.

5.4. Proof of Theorem B. This proof uses the following diagram:

$$D_3 \setminus C \xrightarrow{\beta_1} D_3 \times D_3 \times D_3 \times [3] \xrightarrow{\tilde{\varphi}} D_3 \times D_3 \times D_3 \\ \downarrow (-\partial)^3 \times \text{id} \qquad \qquad \qquad \downarrow (-\partial)^3 \\ \mathbb{N} \times \mathbb{N} \times \mathbb{N} \times [3] \xrightarrow{\varphi} \mathbb{N} \times \mathbb{N} \times \mathbb{N}$$

where notations are:

a)  $\partial^3$  is the mapping defined by

$$-\partial^3(f_1, f_2, f_3) = (-\partial(f_1), -\partial(f_2), -\partial(f_3))$$

b)  $\partial^3 \times \text{id}$  is defined by:

$$-\partial^3 \times \text{id}(f_1, f_2, f_3, i) = (-\partial(f_1), -\partial(f_2), -\partial(f_3), i)$$

c)  $\tilde{\varphi}$  is the mapping given by:

$$\tilde{\varphi}(f_1, f_2, f_3, 1) = (f_2, f_3, f_1), \varphi(f_1, f_2, f_3, 2) = (f_3, f_2, f_1)$$

$$\varphi(f_1, f_2, f_3, 3) = (f_3, f_1, f_2).$$

The proof is then divided in four steps:

1. The diagram commutes as an immediate consequence of the definition of  $\varphi$  and  $\tilde{\varphi}$

2.  $\tilde{\varphi}$  is from  $\tilde{U}$  onto  $D_3 \times D_3 \times D_3$ :

Let  $(f_1, f_2, f_3)$  be an element of  $(D_3)^3$ .  $(-\partial f_1, -\partial f_2, -\partial f_3)$  is in  $\mathbb{N}^3$  and from Lemma 5.4 there exists  $(n_{i_1}, n_{i_2}, n_{i_3}, i)$  in  $U$  such that:

$$i_1, i_2, i_3 \text{ is a permutation of } (1, 2, 3)$$

$$n_i = -\partial f_i \text{ for } i = 1, 2, 3$$

$$\varphi(n_{i_1}, n_{i_2}, n_{i_3}, i) = (-\partial f_1, -\partial f_2, -\partial f_3).$$

We then easily verify that the element  $(f_{i_1}, f_{i_2}, f_{i_3}, i)$  of  $\tilde{U}$  has  $(f_1, f_2, f_3)$  as image by  $\tilde{\varphi}$ .

3.  $\tilde{\varphi}$  is one-to-one: Let  $F = (f_1, f_2, f_3, i), G = (g_1, g_2, g_3, j)$  be such that  $\tilde{\varphi}(F) = \tilde{\varphi}(G)$ . Then  $\partial^3(\tilde{\varphi}(F)) = \partial^3(\tilde{\varphi}(G))$ , and as the diagram commutes we obtain:

$$\varphi(\partial^3 \times \text{id}(F)) = \varphi(\partial^3 \times \text{id}(G)).$$

Now  $\varphi$  is a bijection, and so  $i = j$  and

$$\partial^3(f_1, f_2, f_3) = \partial^3(g_1, g_2, g_3).$$

By the construction of  $\tilde{\varphi}$  this implies  $F = G$ .

4. Theorem B is now established using for  $\beta$  the composition of  $\beta_1$  (a one-to-one mapping from  $D_3 \setminus C$  onto  $\tilde{U}$ ) and  $\tilde{\varphi}$  (a one-to-one mapping from  $\tilde{U}$  onto  $D_3 \times D_3 \times D_3$ ).

REFERENCES

1. P. Cartier and D. Foata, *Problèmes combinatoires de commutation et réarrangement*, Lecture Notes in Math. 55 (Springer-Verlag, Berlin, 1969).
2. R. Cori, *Un code pour les graphes planaires et ses applications*, Astérisque, Société Math. de France 27 (1975).
3. R. Cori and J. Richard, *Énumération des graphes planaires à l'aide de séries formelles en variables non commutatives*, Discrete Math. 2 (1972), 115–162.
4. N. G. de Bruijn and B. J. M. Morselt, *A note on plane trees*, J. Combinatorial Theory 2 (1967), 27–34.

5. J. Edmonds, *Combinatorial representation for polyhedral surfaces*, Notices Amer. Math. Soc. 7 (1960), 646.
6. D. Foata and M. P. Schützenberger, *Théorie géométrique des polynômes eulériens*, Lecture Notes in Math. 138 (Springer-Verlag, Berlin, 1970).
7. A. Jacques, *Sur le genre d'une paire de substitutions*, C. R. Acad. Sci. Paris 267 (1968), 625–627.
8. ——— *Constellations et graphes topologiques* in *Combinatorial theory and its applications*, Colloq. Math. Soc. Janos Bolyai (North Holland Amsterdam, 1970), 657–672.
9. D. A. Klarner, *Correspondance between plane tree and binary sequences*, J. Comb. Theory 9 (1970), 401–411.
10. A. B. Lehman, *A bijective census of rooted planar maps*, Communication at Ontario Math. Conference (1970), unpublished.
11. G. N. Raney, *Functional composition patterns and power series reversion*, Trans. Amer. Math. Soc. 94 (1960), 441–451.
12. R. C. Read, *The coding of various kinds of unlabeled trees* in *Graph theory and computing* (Academic Press, New York, 1972), 153–182.
13. W. T. Tutte, *A census of slicings*, Can. J. Math. 14 (1962), 708–722.
14. ——— *A census of planar maps*, Can. J. Math. 15 (1963), 249–271.
15. ——— *On the enumeration of planar maps*, Bull. Amer. Math. Soc. 74 (1968), 64–74.
16. T. S. Walsh and A. B. Lehman, *Counting rooted maps by genus I and II*, J. Comb. Theory 13B (1972), 192–218 and 122–141.

*Université de Bordeaux I,  
Talence, France*