



# On the Dual König Property of the Order-interval Hypergraph of Two Classes of $N$ -free Posets

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*Abstract.* Let  $P$  be a finite  $N$ -free poset. We consider the hypergraph  $\mathcal{H}(P)$  whose vertices are the elements of  $P$  and whose edges are the maximal intervals of  $P$ . We study the dual König property of  $\mathcal{H}(P)$  in two subclasses of  $N$ -free class.

## 1 Introduction

Let  $(P, \leq)$  be a finite partially ordered set (briefly *poset*  $P$ ). A subset of  $P$  is called a *chain* (resp. *antichain*) if every two elements in  $P$  are comparable (resp. incomparable). The number of elements in a chain is the *length* of the chain. The *height* of an element  $x \in P$ , denoted by  $h(x)$ , is the length of a longest chain in  $P$  having  $x$  as its maximum element. The *height* of a poset  $P$ , denoted  $h(P)$ , is the length of a longest chain in  $P$ . The  *$i$ -level* or *height- $i$ -set* of  $P$ , denoted by  $N_i$ , is the set of all elements of  $P$  that have height  $i$ .

Let  $p$  and  $q$  be two elements of  $P$ . We say  $q$  covers  $p$ , and we denote  $p < q$ , if  $p < v \leq q$  implies  $v = p$ . Furthermore we denote by  $\text{Max } P$  (resp.  $\text{Min } P$ ) the set of all maximal (resp. minimal) elements of  $P$ . A subset  $I$  of  $P$  of the form  $I = \{v \in P, p \leq v \leq q\}$  (denoted  $[p, q]$ ) is called an *interval*. It is maximal if  $p$  (resp.  $q$ ) is a minimal (resp. maximal) element of  $P$ . Denote by  $\mathcal{J}(P)$  the family of maximal intervals of  $P$ . The hypergraph  $\mathcal{H}(P) = (P, \mathcal{J}(P))$  whose vertices are the elements of  $P$  and whose edges are the maximal intervals of  $P$  is said to be the *order-interval hypergraph* of  $P$ .

A subset  $A$  (resp.  $T$ ) of  $P$  is called *independent* (resp. a *point cover* or *transversal set*) if every edge of  $\mathcal{H}$  contains at most one point of  $A$  (resp. at least one point of  $T$ ). A subset  $\mathcal{M}$  (resp.  $\mathcal{R}$ ) of  $\mathcal{J}$  is called a *matching* (resp. an *edge cover*) if every point of  $P$  is contained in at most one member of  $\mathcal{M}$  (resp. at least one member of member of  $\mathcal{R}$ ). Let

$$\begin{aligned}\alpha(\mathcal{H}) &= \max\{|A| : A \text{ is independent}\}, \\ \tau(\mathcal{H}) &= \min\{|T| : T \text{ is a point cover}\}, \\ \nu(\mathcal{H}) &= \max\{|\mathcal{M}| : \mathcal{M} \text{ is a matching}\}, \\ \rho(\mathcal{H}) &= \min\{|\mathcal{R}| : \mathcal{R} \text{ is an edge cover}\}.\end{aligned}$$

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These numbers are called the *independence number*, the *point covering number*, the *matching number*, and the *edge covering number* of  $\mathcal{H}(P)$ , respectively. It is easy to see that  $\nu(\mathcal{H}) \leq \tau(\mathcal{H})$  and  $\alpha(\mathcal{H}) \leq \rho(\mathcal{H})$ . We say that  $\mathcal{H}$  has the *König property* if  $\nu(\mathcal{M}) = \tau(\mathcal{M})$  and *dual König property* if  $\nu(\mathcal{H}^*) = \tau(\mathcal{H}^*)$ , i.e.,  $\alpha(\mathcal{H}) = \rho(\mathcal{H})$ , since  $\alpha(\mathcal{H}) = \nu(\mathcal{H}^*)$  and  $\rho(\mathcal{H}) = \tau(\mathcal{H}^*)$ . This class of hypergraphs has been studied intensively in the past, and we find interesting results from an algorithmic point of view as well as min-max relations [2–8].

Let  $P_1 = (E_1, \leq_1)$  and  $P_2 = (E_2, \leq_2)$  be two posets such that  $E_1$  and  $E_2$  are disjoint. The *disjoint sum*  $P_1 + P_2$  of  $P_1$  and  $P_2$  is the poset defined on  $E_1 \cup E_2$  such that  $x \leq y$  in  $P_1 + P_2$  if and only if  $(x, y \in P_1 \text{ and } x \leq_1 y)$  or  $(x, y \in P_2 \text{ and } x \leq_2 y)$ . The *linear sum*  $P_1 \oplus P_2$  of  $P_1$  and  $P_2$  is the poset defined on  $E_1 \cup E_2$  such that  $x \leq y$  in  $P_1 \oplus P_2$  if and only if  $(x, y \in P_1 \text{ and } x \leq_1 y)$  or  $(x, y \in P_2 \text{ and } x \leq_2 y)$  or  $(x \in P_1 \text{ and } y \in P_2)$ .

Let  $A \subseteq \text{Max } P_1$  and  $B \subseteq \text{Min } P_2$  with  $A$  and  $B$  are not empty. The *quasi-series composition* of  $P_1$  and  $P_2$  denoted  $P = (P_1, A) * (P_2, B)$  is the poset  $P = (E_1 \cup E_2, \leq)$  such that  $x \leq y$  if  $(x, y \in E_1 \text{ and } x \leq_1 y)$  or  $(x, y \in E_2 \text{ and } x \leq_2 y)$  or  $(x \in E_1, y \in E_2)$ , and there exist  $\alpha \in A, \beta \in B$  such that  $x \leq_1 \alpha$  and  $\beta \leq_2 y$ .

## 2 N-free Poset

A poset  $P$  is said to be *series-parallel*, if it can be constructed from singletons  $P_0$  ( $P_0$  is the poset having only one element) using only the two operations disjoint sum and linear sum. It may be characterized by the fact that it does not contain the poset  $N$  as an induced subposet [12, 13].  $P$  is called *N-free* if and only if its Hasse diagram does not contain four vertices  $v_1, v_2, v_3, v_4$ , where  $v_1 < v_2, v_2 > v_3$  and  $v_3 < v_4$ , and  $v_1$  and  $v_4, v_1$  and  $v_3, v_2$  and  $v_4$ , are incomparable. The class of  $N$ -free posets contains the class of series-parallel posets. Habib and Jegou [10] defined the *Quasi-Series-Parallel* (QSP) class of posets, as the smallest class of posets that contains  $P_0$  and closed under quasi-series composition and linear sum. They proved that a poset is  $N$ -free if and only if it is a QSP poset. The following theorem gives many other characterizations of  $N$ -free posets (see [9–11]).

**Theorem 2.1** *The four following properties are equivalent:*

- (i)  $P$  is QSP.
- (ii)  $P$  is an  $N$ -free poset.
- (iii)  $P$  is a C.A.C. (Chain-Antichain Complete) order i.e., every maximal chain intersects each maximal antichains.
- (iv) The Hasse diagram of  $P$  is a line-digraph.
- (v) For every two elements  $p, q \in P$ , if  $N(p) \cap N(q) \neq \emptyset$ , then  $N(p) = N(q)$ , where  $N(p)$  denoted the set of all elements of  $P$  that cover  $p$  in  $P$ .

It is known that the order-interval hypergraph  $\mathcal{H}(P)$  has the König and dual König properties for the class of series-parallel posets [3]. In [4], it was proved that  $\mathcal{H}(P)$  has again the dual König property for the class of a posets that contains the series-parallel posets and whose members have comparability graphs that are distance-hereditary graphs or generalizations of them. If  $P$  is an  $N$ -free poset, the König property is not satisfied in general; see [4]. The poset of Figure 1 is an example where  $\nu(\mathcal{H}(P)) = 1$ ,

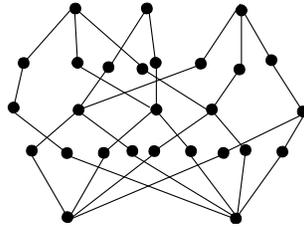


Figure 1:  $v(\mathcal{H}(P)) = 1$  and  $\tau(\mathcal{H}(P)) = 2$

$\tau(\mathcal{H}(P)) = 2$ . In this paper, we consider two classes of N-free posets and prove that the dual König property of the order-interval hypergraph of these classes of posets are satisfied.

### 2.1 Blocks in an N-free Poset

There is a useful representation of an N-free poset, namely the *block* (see [1]). If  $P$  is an N-free poset with levels  $N_1, \dots, N_r$ , a block of  $P$  is maximal complete bipartite graph in the Hasse diagram of  $P$ . More precisely, a block of  $P$  is a pair  $(A_i, B_i)$ , where  $A_i, B_i \subset P$  such that  $A_i$  is the set of all lower covers of every  $x \in B_i$  and  $B_i$  is the set of all upper covers of every  $y \in A_i$ . By convention,  $(\emptyset, \text{Min } P)$  and  $(\text{Max } P, \emptyset)$  are blocks

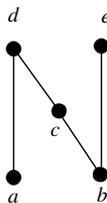


Figure 2:  $P$  is N-free with blocks  $(\emptyset, \{a, b\})$ ,  $(\{b\}, \{c, e\})$ ,  $(\{a, c\}, \{d\})$  and  $(\{d, e\}, \emptyset)$ .

In this paper, we say that  $(A_i, B_i)$  and  $(A_j, B_j)$  are *adjacent* if there exists at least one vertex of  $A_i \cup B_i$  in the same interval in  $P$  with at least one vertex of  $A_j \cup B_j$ . For example, the blocks  $(\{b\}, \{c, e\})$  and  $(\{a, c\}, \{d\})$  of poset of Figure 2 are adjacent.

### 2.2 N-free Poset of Type 1

**Definition 2.2** Let  $P$  be a connected poset with levels  $N_1, N_2, \dots, N_r$ . We say that  $P$  is of *Type 1* if there exists an integer  $n$  such that the induced subposet  $P_{n, n+1}$  formed from the consecutive levels  $N_n \cup N_{n+1}$  is of the form  $N_n \oplus N_{n+1}$ .

For the class of posets of Type 1, we give the following result.

**Theorem 2.3** Let  $P$  be a poset of Type 1. Then  $\mathcal{H}(P)$  has the dual König property, and we have  $\alpha(\mathcal{H}(P)) = \rho(\mathcal{H}(P)) = \text{Max}\{|\text{Max } P|, |\text{Min } P|\}$ .

**Proof** We set  $\text{Min } P = \{a_1, a_2, \dots, a_k\}$  and  $\text{Max } P = \{b_1, b_2, \dots, b_l\}$ . Consider the family of edges  $\mathcal{J}$  of  $\mathcal{H}(P)$  such that

$$\mathcal{J} = \begin{cases} \{[a_j, b_j], j = 1, \dots, k\} \cup \{[a_k, b_j], j = k + 1, \dots, l\} & \text{if } k \leq l, \\ \{[a_j, b_j], j = 1, \dots, l\} \cup \{[a_j, b_l], j = l + 1, \dots, k\} & \text{if } k > l. \end{cases}$$

It is not difficult to see that  $\mathcal{J}$  is an edge-covering family of  $\mathcal{H}(P)$  of cardinality equal to  $\text{Max}\{|\text{Max } P|, |\text{Min } P|\}$ . Hence,  $\alpha(\mathcal{H}(P)) = \rho(\mathcal{H}(P)) = \text{Max}\{|\text{Max } P|, |\text{Min } P|\}$  ■

In particular, the order-interval hypergraph of the N-free poset of Type 1 has the dual König property.

### 3 N-free Poset of Type 2

#### Definitions

(a) Let  $P$  be a connected N-free poset with levels  $N_1, N_2, \dots, N_r$ . We say that  $P$  is a poset of Type 2 if there exists an integer  $n$  such that  $N_n$  is the first level where the induced subposet  $P_{n,r}$  is disconnected of the form  $P_{n,r} = P_1 + P_2 + \dots + P_l$ , and for all  $i \in L = \{1, \dots, l\}$ ,  $P_i$  is connected poset of Type 1.

(b) We say that the subposet  $P_i$  is linked with the subposet  $P_j$  by a vertex  $z$  of  $N_1$ , if we can obtain intervals of the form  $[z, x]$  and  $[z, y]$  for each  $x \in \text{Max } P_i$  and  $y \in \text{Max } P_j$ , and we say  $z$  links  $P_i$  with  $P_j$ .

(c) We say that  $P_i$  is linked with  $P_j$  by the subset  $R$  of  $N_1$  if for every element  $z$  of  $R$ ,  $z$  links  $P_i$  with  $P_j$ .

**Example 3.1** The poset  $P$  in Figure 3 is N-free of Type 2; it is easy to see that  $N_2$  is the first level where  $P_{2,3} = P_1 + P_2$  is disconnected poset where  $P_1$  and  $P_2$  are posets of Type 1. On the other hand,  $Q$  is an N-free poset but not of Type 2.

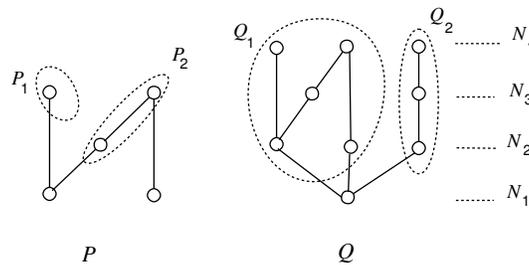


Figure 3:

In order to prove the dual König property of  $\mathcal{H}(P)$ , where  $P$  is N-free of Type 2, let us introduce the following notation.

**Notation**

(a) For every subposet  $P_k$ , we denote by  $R_k$  the subset of  $N_1$ , where every element of  $R_k$  is comparable with all elements of  $\text{Max } P_k$ , and  $R_k$  does not link  $P_k$  with any other poset  $P_s, s \in L$ . The set  $R_k$  can be empty.

(b) For every subposet  $P_k$ , we denote by  $R'_{ik}, i \in I_k = \{1, 2, \dots, |N_1|\}$ , the subset of  $N_1$  that links  $P_k$  with the same family of poset  $\{P_s\}_{s \in L}$ . We can obtain  $R'_{ik} = R'_{jl}$  for  $i \neq j$  and  $k \neq l$ .

**Observation 3.2** *The family  $\{R'_{ik}\}_{k \in L, i \in I_k}$  is pairwise disjoint.*

See Figure 4 for an illustration of the class of N-free posets of Type 2.

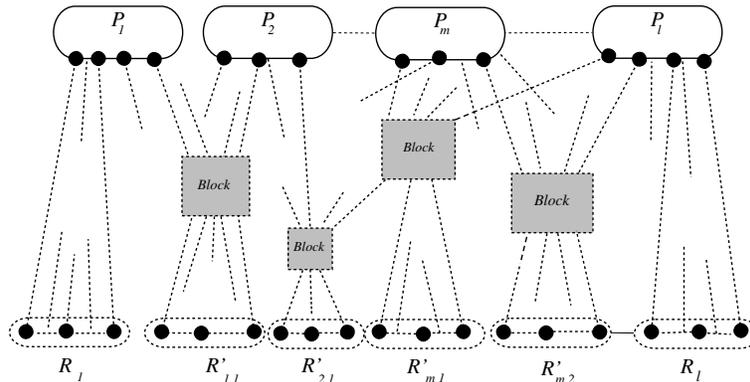


Figure 4: Illustration of an N-free poset of Type 2

**3.1 Maximal Stable Sets of  $\mathcal{H}(P)$**

In our poset, it is clear that for a linked subposet family  $F_k = \{P_l\}_{l \in L}$ , we can obtain blocks  $(A_i, B_i)$  in the level  $N_{n-j}$ , for  $j \in \{0, 1, \dots, n-1\}$ , i.e.,  $B_i$  intersects  $N_{n-j}$ , and every element  $x$  of  $A_i$ ,  $x$  links a subfamily  $F_s$  of  $F_k$ , we say  $(A_i, B_i)$  links  $F_s$ . Such blocks must exist in  $N_n$  since  $P$  is N-free poset of Type 2.

We note the following observation.

**Observation 3.3** *For every block  $(A_i, B_i)$  that links  $F_s$ ,  $B_i$  has the following partition:  $B_i = \cup_{t \in T} B_{i,t}$ , where  $\forall x \in B_{i,t}$ ,  $x$  is comparable with a vertex of  $\text{Min } P_t$ , where  $P_t \in F_s$ , and  $|F_s| = |T|$*

Let us now give two algorithms to find maximal stable sets of an N-free poset of Type 2; the second algorithm can be applied only after the first.

**Maximal Stable-set 1 Algorithm**

INPUT: An N-free poset  $P$  of Type 2.  $F_1, F_2, \dots, F_m$  all linked subposet families of  $P$ .

- (a) For each  $k$ , from  $k = 1$  to  $m$ .
- (b) For each  $j$ , from  $j = 0$  to  $n - 1$ , in  $N_{n-j}$  we determine  $C_{k,j}$  by taking for every block  $(A_i, B_i)$  that links a subfamily of  $F_k$ , one vertex from each  $B_{i,t}$  such that:
  - (i) if there exists a family  $\{B_{i,t}\}_i$  from block family that are adjacent pairwise, we take only one vertex from only one set of  $\{B_{i,t}\}_i$ ;
  - (ii) we delete every vertex which is in the same interval with a vertex of  $C_{k,t}$ ,  $t < j$ .
- (c) Put  $C_k = \bigcup_{j=0}^{n-1} C_{k,j}$ .
- (d) Output  $\mathcal{C} = (\bigcup_{k=1}^m C_k) \cup (\bigcup_{l \in L} R_l)$ . End

**Theorem 3.4** *The set  $\mathcal{C}$  is a maximal stable set of  $\mathcal{H}(P)$ .*

**Proof**  $\mathcal{C}$  is a stable set by construction of every  $C_k$ . It remains the maximality of  $\mathcal{C}$ . We say that an interval  $I$  crosses a block  $(A_i, B_i)$  if  $I$  intersects  $B_i$ . Let us show that for every interval  $I$  of  $P$ ,  $I$  contains one vertex of  $\mathcal{C}$ , and this means that for every  $x \in P$ ,  $\mathcal{C} \cup \{x\}$  will not be a stable set.

In the case where  $I$  does not cross any block, the minimal vertex of  $I$  will be in  $R_l$ . Now, in the case where  $I$  crosses a block  $(A_i, B_i)$ , let  $y$  be a commun vertex of  $B_i$  and  $I$ . If  $y \in \mathcal{C}$ , then  $I$  intersects  $\mathcal{C}$ . Otherwise,  $y \notin \mathcal{C}$ , which means that  $y$  is in the same interval  $J$  with an element  $y'$  of  $\mathcal{C}$ . Consequently,  $I$  and  $J$  will have minimal vertices in  $R'_{pq}$  and maximal vertices in  $\text{Max } P_i$ ; this gives  $y' \in I$ . ■

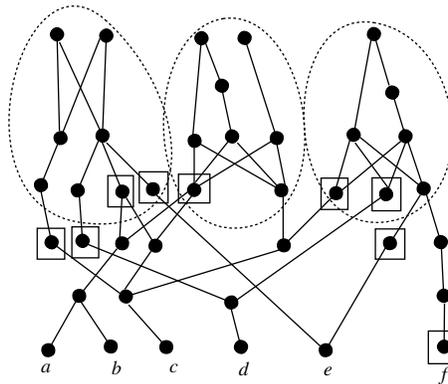


Figure 5: An N-free poset  $P$  of Type 2. Applying the Maximal Stable-Set 1 algorithm on  $P$ ; the framed vertices form a maximal stable set of  $\mathcal{H}(P)$ .

**Example 3.5** The poset of Figure 5 is N-free of Type 2, where  $P_1, P_2$ , and  $P_3$  are the subposets surrounded from left to right. We have  $R'_{11} = R'_{12} = \{a, b\}$ ,  $R'_{21} = R'_{22} =$

$R'_{13} = \{c\}$ ,  $R'_{31} = R'_{23} = \{d\}$ ,  $R'_{41} = R'_{33} = \{e\}$ , and  $R_3 = \{f\}$ . The framed vertices form the maximal stable set  $\mathcal{C}$  of  $\mathcal{H}(P)$  obtained by the Maximal Stable-set 1 algorithm.

We will need the following definition.

**Definition 3.6** In  $\mathcal{H}(P)$ , for every vertex  $x \in P$ , a *stable adjacent*  $M_x$  to  $x$  is the set of all vertices  $y$  such that  $x$  and  $y$  are in the same interval of  $P$ , where  $M_x$  is stable.  $M_x$  can be equal to  $\{x\}$ . We say  $M_D$  is a stable adjacent to the set  $D$  of  $P$  if it is a maximal stable subset of the union of all  $M_x$ ,  $x \in D$ , where  $|M_D| \geq |D|$ .

We can write  $\bigcup_{k=1}^m C_k = D_1 \cup D_2 \cup \dots \cup D_m$  for the stable set obtained from the Maximal Stable-set 1 algorithm, where  $D_i$  are subsets of blocks of  $P$ . We determine a new maximal stable set  $\mathcal{C}'$  from  $\mathcal{C}$  as follows.

**Maximal Stable-set 2 Algorithm**

INPUT: An N-free poset  $P$  of Type 2, and maximal stable set  $\mathcal{C}$ .

OUTPUT: A new maximal stable set  $\mathcal{C}'$ .

1.  $\mathcal{C}' := \mathcal{C}$ .
2. For each  $i$ , from  $i = 1$  to  $m$ .
3. We determine  $M_{D_i}$  the stable adjacent to  $D_i$  such that  $\mathcal{C} - (\bigcup_{t=1}^{t=i} D_t) \cup (\bigcup_{t=1}^{t=i} M_{D_t})$  is stable.
4. We take  $\mathcal{C}' := \mathcal{C} - (\bigcup_{t=1}^{t=i} D_t) \cup (\bigcup_{t=1}^{t=i} M_{D_t})$ .
5. Stop.

By construction of  $\mathcal{C}'$ , we deduce the following result.

**Proposition 3.7** The set  $\mathcal{C}'$  is a maximal stable set of  $\mathcal{H}(P)$ .

We denote by  $C'_k$  the set of all vertices obtained from every  $x_i \in C_k$  using the Maximal Stable-set 2 algorithm.

As a consequence of the previous algorithms, we make the following observation.

**Observation 3.8** Consider the subposet family  $F_k$  linked by  $R'_{pq}$ .

(i) The set  $R'_{pq}$  has the following partition:  $R'_{pq} = \bigcup_s R'_{pq,s}$ , where for every  $s$ ,  $R'_{pq,s}$  is a stable adjacent to  $A_s$  a subset of  $C'_k$ .

(ii) It will be possible to obtain that the family  $\{A_s\}_s$  is pairwise disjoint.

**Proof** To prove the second observation, we suppose that  $x$  is a common vertex of  $A_s$  and  $A_{s'}$ . Let  $I$  (resp.  $J$ ) an interval containing  $x$  with minimal element  $c_j \in R'_{pq,s}$  (resp.  $c_{j'} \in R'_{p'q',s'}$ ). In  $I$  (resp.  $J$ ) there exists a vertex  $z$  (resp.  $z'$ ) that is incomparable with every vertex of  $R'_{p'q',s'}$  (resp.  $R'_{pq,s}$ ) (we take as an example, the vertex  $z$  (resp.  $z'$ ) such that  $c_j < z$  (resp.  $c_{j'} < z'$ )). Otherwise, we will obtain  $R'_{pq,s} = R'_{p'q',s'}$  since  $P$  is N-free. In this case, we can reconstruct  $\mathcal{C}$  by starting with  $z$  and  $z'$  respectively to obtain two new disjoint sets. ■

In the remainder of this paper, we suppose that  $\mathcal{C}'$  verifies Observation 3.8(ii).

**Example 3.9** The poset of Figure 6 is N-free of Type 2, where  $\mathcal{C} = \{a, b\}$ . Applying the Maximal Stable-set 2 algorithm, we obtain two different maximal stable sets:  $\mathcal{C}'_1$  is the framed vertex set and  $\mathcal{C}'_2$  is the surrounded vertex set. We remark that  $\mathcal{C}'_2$  verifies Observation 3.8(ii), while  $\mathcal{C}'_1$  does not.

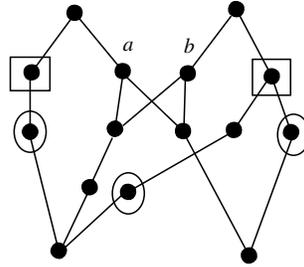


Figure 6: Two different maximal stable sets of  $\mathcal{H}(P)$  by applying the Maximal Stable-set 2 algorithm.

### 3.2 Edge Covering Family of $\mathcal{H}(P)$

In this section, we will present an algorithm to construct an edge covering family of  $\mathcal{H}(P)$  where  $P$  is an N-free of Type 2.

We set

$$\begin{aligned} \text{Max } P_l &= \{b_1^l, b_2^l, \dots, b_{|\text{Max } P_l}^l\}, & R_l &= \{a_1, a_2, \dots, a_{|R_l|}\}, \\ R'_{pq,s} &= \{c_1, c_2, \dots, c_{|R'_{pq,s}|}\}, & \bigcup_{i \in I_l} R'_{il} &= \{c'_1, c'_2, \dots, c'_{m_l}\}. \end{aligned}$$

**Theorem 3.10** *If for every  $k \in L$  we have*

$$(3.1) \quad |\text{Max } P_k| \geq |R_k| + \sum_{i \in I_k} |R'_{ik}|,$$

*then  $\mathcal{H}(P)$  has the dual König property and  $\alpha(\mathcal{H}(P)) = \rho(\mathcal{H}(P)) = |\text{Max } P|$ .*

**Proof** For every  $P_k$ , we consider the edge family:

$$\begin{aligned} \mathcal{J}_k &= \{[a_i, b_i], i = 1, \dots, |R_k|\} \cup \{[c'_{j-|R_k|}, b_j], j = |R_k| + 1, \dots, |R_k| + m_k\} \\ &\quad \cup \{[c'_{m_k}, b_s], s = m_k + |R_k| + 1, \dots, |\text{Max } P_k|\}. \end{aligned}$$

The union of all  $\mathcal{J}_k, k \in L$  is an edge covering family of  $\mathcal{H}(P)$  with cardinality equal to  $|\text{Max } P|$  and as  $\text{Max } P$  is a stable set of  $\mathcal{H}(P)$  then  $\alpha(\mathcal{H}(P)) = \rho(\mathcal{H}(P)) = |\text{Max } P|$ . ■

We remark that by applying the Maximal Stable-set 2 algorithm to  $P$ , we can obtain different maximal stable sets of  $\mathcal{H}(P)$ , and this depends on the choice of  $M_{D_i}$ . In the next algorithm we need to characterize the set  $\mathcal{C}'$  as follows.

$\mathcal{C}'$  is determined such that for every subposet family  $F_k$  that contains subposets  $P_l$  verifying (3.1), we determine  $M_{D_i}$  different to  $D_i$ , and if  $x \in D_i$  is incomparable with all vertices of  $\text{Max } P_l$ , then  $M_x$  will be too. For other subposet families,  $M_{D_i}$  does not contain a vertex of  $\text{Max } P_m$ , where  $R_m$  is not empty.

**Edge-Cover Algorithm**

INPUT: An N-free poset  $P$  of Type 2 and the maximal stable set  $\mathcal{C}'$ .

OUTPUT: An edge covering family  $\mathcal{J}(\mathcal{H}(P))$ .

**Step 1** For every  $R_l$ , where  $P_l$  does not verify (3.1), we construct the edge family  $E_l$  with  $|R_l|$  intervals as follows:

1.1 If  $|R_l| \leq |\text{Max } P_l|$ :  $E_l = \{[a_j, b_j^l], j = 1, 2, \dots, |R_l|\}$ .

1.2 Otherwise:  $E_l = \{[a_j, b_j^l], j = 1, 2, \dots, |\text{Max } P_l|\} \cup \{[a_t, b_{|\text{Max } P_l}|], t = |\text{Max } P_l| + 1, \dots, |R_l|\}$ .

**Step 2** For every  $P_l$ , where  $P_l$  verifies (3.1), we construct the edge family  $J_l$  as follows:

$$J_l = \{[a_i, b_i], i = 1, \dots, |R_l|\} \cup \{[c'_{j-|R_l|}, b_j], j = |R_l| + 1, \dots, |R_l| + m_l\} \\ \cup \{[c'_{m_l}, b_s], s = m_l + |R_l| + 1, \dots, |\text{Max } P_l|\}.$$

We obtain  $|\text{Max } P_l|$  intervals.

**Step 3** First, determine all linked subposet families  $F_1, F_2, \dots, F_m$ . Then apply this step to  $F_k = \{P_l\}_{l \in S_k}$ , which is linked by  $R'_{pq}$  for  $k = 1$  to  $k = m$ .

In this step, we use the vertices  $b_t^l$  of  $\text{Max } P_l$ ,  $P_l \in F_k$ , which are not used in Step 1 or in the application of this step to  $F_t$ , where  $t < k$ ; otherwise, we use vertices already used.

Let  $A'_s$  be the set  $A_s$  deleting all vertices comparable with  $\text{Max } P_m$ , where  $P_m$  verifies (3.1), and  $F'_k = \{P_l\}_{l \in S'_k}$  be the family  $F_k$  deleting all subposets verifying (3.1). For every  $R'_{pq,s}$  we construct the edge family  $I_s$  as follows:

3.1 If  $|A'_s| \leq |R'_{pq,s}|$ , then  $I_s = \{[c_j, b_t^l], j = 1, 2, \dots, |A'_s| \text{ and } l \in S'_k\}$ . We obtain  $|A'_s|$  intervals.

3.2 If  $|A'_s| > |R'_{pq,s}|$ , then

$$I_s = \{[c_j, b_t^l], j = 1, 2, \dots, |R'_{pq,s}| \text{ and } l \in S''_k \subset S'_k\} \cup \{[c_l, b_t^l], l \in (S'_k - S''_k)\}.$$

We obtain  $|A'_s|$  intervals.

**Step 4** It remains some minimal vertices  $c_j$  that are not used in Steps 1 and 3 such that  $c_j \in R'_{pq,s}$  and  $R'_{pq}$  does not link any subposet verifying (3.1). In this step, we construct  $J_{c_j}$  the interval containing  $c_j$  and  $b_t^l$  a maximal vertex that is not already used, otherwise,  $J_{c_j}$  is any interval containing  $c_j$ .

**Step 5** We take  $\mathcal{J}(\mathcal{H}(P))$  to be the set of all intervals obtained from Step 1 to Step 4.

END

**Theorem 3.11** *The Edge-Cover algorithm applied to an N-free poset  $P$  of Type 2, yields an edge-covering family of  $\mathcal{H}(P)$ .*

**Proof** We can assert that every  $z$  of  $P$  that is a minimal element, comparable with a vertex of  $R_m$  or comparable with a vertex of  $\text{Max } P_l$ , where  $P_l$  verifies (3.1) is covered by  $\mathcal{J}(\mathcal{H}(P))$ .

Moreover, if  $z > x$ , where  $x \in A'_s$ , then  $z$  would be covered by the interval of  $\mathcal{J}(\mathcal{H}(P))$  that intersects  $A'_s$ .

In other cases, suppose that there exists  $z$  of  $P$  that is not covered by  $\mathcal{J}(\mathcal{H}(P))$ . We distinguish two cases.

*Case 1.* If  $z$  is a maximal of  $P_l$  and no interval obtained from Step 3 or Step 4 covers  $z$ , then  $P_l$  necessarily would verify (3.1). This contradicts the construction of intervals in these steps.

*Case 2.* Let  $J \notin \mathcal{J}(\mathcal{H}(P))$  containing  $z$  and  $x$ , where  $x \in A'_s$  and  $x \not\leq z$ . Let  $I$  be the interval of  $\mathcal{J}(\mathcal{H}(P))$  containing  $x$ . The only form of  $I$  and  $J$  is that they will have maximal elements in  $\text{Max } P_l$  and two different minimal elements in  $R'_{pq,s}$ . Then  $z$  is not covered by  $I$ , then for every couple  $(t, t')$  of  $(I, J)$ , where  $t \leq x$  and  $t' \leq z$ , we will have  $t \not\leq t'$ . We suppose that such a couple exists.

If  $t$  and  $t'$  are not in the same interval and  $A'_s \cup \{t, t'\} - \{x\}$  is stable, then  $x$  can be replaced by  $t$  and  $t'$  in  $\mathcal{C}'$ , and this contradicts the construction of  $\mathcal{C}'$ . Otherwise, we can reconstruct  $A'_s$  starting by  $z$ . In this case,  $R'_{pq,s}$  will be partitioned into at least two subsets, and by applying the Edge-Cover algorithm;  $z$  will be covered by the new family. ■

As a consequence of Theorem 3.11, we have the following corollary.

**Corollary 3.12** *If in the Edge-Cover algorithm, for every vertex  $x$  of  $\text{Max } P$  (resp.  $\text{Min } P$ ),  $x$  is taken only once in the construction of  $\mathcal{J}(\mathcal{H}(P))$ , then  $P$  will have the dual König property.*

**Proof** In this case, we will have  $|\mathcal{J}(\mathcal{H}(P))| = |\text{Max } P|$  (resp.  $|\text{Min } P|$ ), and as  $\text{Max } P$  and  $\text{Min } P$  are stable sets of  $\mathcal{H}(P)$ , therefore

$$\alpha(\mathcal{H}(P)) = \rho(\mathcal{H}(P)) = |\text{Max } P| \quad (\text{resp.}, \alpha(\mathcal{H}(P)) = \rho(\mathcal{H}(P)) = |\text{Min } P|).$$

■

**Theorem 3.13** *Let  $P$  be an  $N$ -free poset of type 2. Then  $\mathcal{H}(P)$  has the dual König property.*

**Proof** The main idea of the proof is to use  $\mathcal{J}(\mathcal{H}(P))$  obtained from the Edge-Cover algorithm for constructing a stable set  $\mathcal{C}(\mathcal{H})$  of  $\mathcal{H}(P)$  with the same size as  $\mathcal{J}(\mathcal{H}(P))$ .

Let  $B_1$  (resp.  $B_2$ ) be the union of all  $R_l$  (resp.  $\text{Max } P_k$ ), where  $P_l$  (resp.  $P_k$ ) does not verify (resp. verifies) (3.1).

From Step 1 (resp. Step 2) of the Edge-Cover algorithm,  $B_1$  (resp.  $B_2$ ) is a stable set with the cardinality equal to the cardinality of the union of all  $E_l$  (resp.  $J_l$ ). It becomes clear that  $B_1 \cup B_2$  is stable set.

The union of all  $I_s$  of Step 3.1 can be partitioned into 2 subsets. The first denoted by  $D_1$ , which is the union of all  $I_s$ , where  $R'_{pq,s}$  does not link subposets verifying (3.1), and the second is denoted by  $D_2$ . Let  $B_{3,1}$  be the union of all  $R'_{pq,s}$ , where  $R'_{pq}$  does

not link subposets verifying (3.1) and  $|R'_{pq,s}| > |A_s|$ .  $B_{3,1}$  is a stable set with cardinality equal to  $|D_1|$  plus the cardinality of the union of all  $J_{c_j}$  of Step 4.

We denote by  $B_{3,2}$  the union of all  $A'_s$  such that  $|A'_s| > |R'_{pq,s}|$  or  $|A'_s| \leq |R'_{pq,s}|$ , where  $R'_{pq}$  links subposets verifying (3.1). From Observation 3.8(ii), we deduce that there is no common vertex  $x$  of  $A_s$  and  $A_{s'}$  that is covered by two different intervals of  $\mathcal{J}(\mathcal{H}(P))$ . Consequently,  $|B_{3,2}|$  is equal to  $|D_2|$  plus the cardinality of the union of all  $I_s$  of Step 3.2. Consider the set  $\mathcal{C}(\mathcal{H}) = B_1 \cup B_2 \cup B_{3,1} \cup B_{3,2}$ .

Hence, it is not difficult to see that  $\mathcal{C}(\mathcal{H})$  is a stable set with size  $|\mathcal{J}(\mathcal{H}(P))|$ . ■

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