

# A primitive group ring

Warren Brisley and R. Groenhout

An explicit construction is given for a primitive group ring, together with an explicit construction of a faithful irreducible module for it.

Until Formanek and Snider established sufficient conditions [1] for a group to generate a primitive group ring, there were some doubts about the existence of such objects. Their proof of primitivity uses the "internal" characterisation of primitivity in terms of the existence of a certain maximal one-sided ideal. By contrast, it seems worthwhile to construct a particularly easily-described group ring, together with an explicit faithful irreducible module for it; this provides an actual example, in which the primitivity is displayed in its "external" characterisation.

## 1. The group

Let  $\Sigma$  be the group of permutations, on the non-negative integers  $N$ , generated by  $\{\sigma_0, \sigma_1, \sigma_2, \dots\}$ , where  $\sigma_i$  is defined by:

*( $n\sigma_i$ ) is the number obtained from  $n$  by changing the digit for  $2^i$  in the binary representation of  $n$ .*

(Thus,  $\sigma_0$  interchanges each even integer with its successor,  $\sigma_1$  permutes  $N$  to  $2, 3, 0, 1, 6, 7, 4, 5, \dots$ , and in general,  $\sigma_k$  interchanges (rigidly) blocks of length  $2^k$ .) Clearly  $\Sigma$  is singly transitive, any element of  $\Sigma$  is specified completely by its action on

---

Received 6 December 1974. The authors thank R.W. Robinson for the encouragement to believe that a set like  $S$  could be chosen in an easily-described manner.

0, and  $\Sigma$  is isomorphic to  $C_2 \times C_2 \times C_2 \times \dots$ .

Let  $A$  be another copy of  $C_2 \times C_2 \times C_2 \times \dots$  generated by commuting involutions  $\{a_0, a_1, a_2, \dots\}$  and split-extend  $A$  by  $\Sigma$  using the automorphisms  $\sigma_j : a_i \rightsquigarrow a_k$  where  $k = (i\sigma_j)$ . The result is the group  $G = \langle \{a_0, a_1, a_2, \dots\} \cup \{x_0, x_1, x_2, \dots\} \rangle$  with the relations

- (i) all  $x_i$  are of order 2 and commute pairwise,
- (ii) all  $a_i$  are of order 2 and commute pairwise,
- (iii) for each  $i, j$  pair,  $x_j a_i x_j = a_k$  where  $k = (i\sigma_j)$ .

By the transitivity of  $\Sigma$ ,  $G$  is generated by  $\{a_0\} \cup \{x_0, x_1, x_2, \dots\}$ . Further, by the block action of elements of  $\Sigma$ ,  $G$  is locally finite; and by the transitivity of  $\Sigma$ , any normal subgroup of  $G$  must be infinite. (As an aid to calculation, note that any element of  $G$  can be written as  $ax$ , with  $a \in A$ ,  $x \in X$ , where  $X$  is generated by  $\{x_0, x_1, x_2, \dots\}$ .)

## 2. The group ring

Let  $F$  be any field: the group ring  $F(G)$  consists of elements (formal sums) of the form  $\left[ \sum_i \alpha_i g_i \right]$  with  $\alpha_i \in F$ ,  $g_i \in G$ ; addition and multiplication are defined in the natural way, using the multiplication in  $G$  and collecting terms.

In this particular case, we require that  $F$  not have characteristic 2. We note that for this particular group ring, any element  $r$  of  $F(G)$  can be written

$$r = A_1 X_1 + A_2 X_2 + \dots + A_k X_k$$

where

- (i) each  $X_i$  is an element of  $X$ , and all the  $X_i$  are different,
- (ii) each  $A_i$  is an element of  $F(A)$ , and so it can be written

(with  $e$  the unit element of  $G$ ):

$$A_i = \beta e + \beta_0 a_0 + \beta_1 a_1 + \beta_{01} a_0 a_1 + \dots + \beta_{01\dots n} a_0 a_1 \dots a_n$$

where there are  $2^{n+1}$  terms for some  $n$ . (Some, but not all, of the  $\beta$ 's may well be zero: each  $\beta$  is in  $F$ .)

### 3. The module

Let  $V$  be the vector space over  $F$  with basis (independent) elements  $\{b_0, b_1, b_2, \dots\}$ : that is,  $V$  consists of finite formal sums  $\sum \alpha_i b_i$ , with the  $\alpha_i$  in  $F$ . We need only define the action of  $G$  on  $V$  (in fact, on the basis elements) in such a way as to make  $V$  into a  $G$ -module, and then  $V$  will be an  $F(G)$ -module in the natural way.

Assuming the existence of a certain set

$$S = \{2, 5, 6, 7, 12, \dots\}$$

(whose existence, construction and use we will deal with later), we specify:

$$b_i x_j = b_k \quad \text{where } k = (i\sigma_j),$$

$$b_i a_0 = \epsilon_i b_i \quad \text{where } \epsilon_i = -1 \text{ if } i \in S,$$

$$\epsilon_i = +1 \text{ if } i \notin S.$$

This specification extends associatively to words in the generators of  $G$ , and we obtain  $V$  as a  $G$ -module if the relations in  $G$  are satisfied. Clearly the requirements of order and commutativity are satisfied: for the other relations, we note that elements of  $\Sigma$  are uniquely specified by their action on  $0$ , so if  $x_j$  interchanges  $r$  and  $n$ , and  $X_r$  interchanges  $0$  and  $r$  then  $x_j X_r$  is precisely the element which interchanges  $0$  and  $n$ . Thus  $b_i(x_j a_r x_j)$  is  $b_i(x_j X_r a_0 X_r x_j)$ , which is  $b_i((x_j X_r) a_0 (x_j X_r))$ , which is  $b_i a_n$ . In either case, the result is  $\epsilon_t b_i$ , where  $x_j X_r$  interchanges  $i$  and  $t$ .

$V$  is now an  $F(G)$ -module in the natural way.

4. Faithfulness and irreducibility

The set  $S$  is selected using the following array:

0	1	2	3	4	5	6	7	8	9	...
1	0	3	2	5	4	7	6	9	8	...
2	3	0	1	6	7	4	5	10	11	...
3	2	1	0	7	6	5	4	11	10	....
.	.	.	.	.	.	.	.	.	.	.
.	.	.	.	.	.	.	.	.	.	.

where the entry  $e_{ij}$  in row  $i$ , column  $j$  is  $j\sigma$ , where  $\sigma$  is that unique element of  $\Sigma$  which interchanges 0 and  $i$ . (The rows and columns are numbered from zero.) Clearly, if  $e_{ij} = n$ , then the action of  $a_i$  on  $b_j$  is given by:

$$b_j a_i = -b_j \text{ if and only if } b_n a_0 = (-b_n), \text{ that is if } n \in S,$$

$$b_j a_i = +b_j \text{ if and only if } n \text{ is not in } S.$$

The array is symmetric, since by the definition of  $\Sigma$ ,  $e_{ij}$  is calculated from  $j$  using the difference between the binary expression of zero and that of  $i$ ; this is equivalent to using the non-zero binary digits of  $j$  to mark which of the binary digits of  $i$  to change; thus  $e_{ij} = e_{ji}$ .

We now need a lemma:

LEMMA. *The set  $S$  can be chosen in such a way that, given  $n \geq 0$ , the set of part-columns of length  $n + 1$ :*

$$\left\{ \begin{bmatrix} e_{0j} \\ e_{1j} \\ \vdots \\ e_{nj} \end{bmatrix}, j = 0, 1, 2, \dots \right\}$$

*contains at least one of each of the  $2^{n+1}$  possible patterns of +, - produced by the  $\epsilon_k$ , where  $k = e_{ij}$ .*

(Thus, with  $S$  as mentioned in the previous section, and with

$n = 1$  , we have the pattern

$$\left. \begin{array}{l} b_0 a_0 = +b_0 \\ b_0 a_1 = +b_1 \end{array} \right\} \left. \begin{array}{l} b_2 a_0 = -b_2 \\ b_2 a_1 = +b_2 \end{array} \right\} \left. \begin{array}{l} b_4 a_0 = +b_4 \\ b_4 a_1 = -b_4 \end{array} \right\} \left. \begin{array}{l} b_6 a_0 = -b_6 \\ b_6 a_1 = -b_6 \end{array} \right\}$$

corresponding to columns  $j = 0, 2, 4, 6$  , and rows 0 and 1 for  $a_0$  and  $a_1$  acting on those  $b_j$  : for example,  $b_4 a_1 = -b_4$  since  $e_{14} \in S$  .)

Leaving the algorithm to produce  $S$  until later, we now have:

(i)  $V$  is a faithful  $F(G)$ -module.

Proof. Take any  $r$  in  $F(G)$  , and assume  $Vr = 0$  . Write  $r$  as in Section 2. Then

$$b_j r = b_{j1} A_{j1} X_1 + b_{j2} A_{j2} X_2 + \dots + b_{jk} A_{jk} X_k = 0 \text{ for each } j .$$

This reads:

$$\alpha_{1,j} b_j X_1 + \alpha_{2,j} b_j X_2 + \dots + \alpha_{k,j} b_j X_k = 0 ,$$

where the  $\alpha_{i,j}$  are in  $F$  . Since the  $X_1, \dots, X_k$  are all different, so too are the  $b_j X_1, \dots, b_j X_k$  , by the properties of  $\Sigma$  . Hence each  $\alpha_{i,j}$  is zero. Now any  $\alpha_{i,j}$  is produced from  $A_i$  ( $b_j A_i = \alpha_{i,j} b_j$ ) by

$$A_i = \beta e + \beta_0 a_0 + \beta_1 a_1 + \dots + \beta_{01\dots n} a_0 a_1 \dots a_n$$

for some  $n$  . Then

$$\alpha_{i,j} = \beta \pm \beta_0 \pm \beta_1 \pm \dots \pm \beta_{01\dots n}$$

where the  $\pm$  signs depend on the allotment of the  $\epsilon_k$  . Since each of the  $2^{n+1}$  patterns of  $\pm$  which could be produced by the action of the  $a_i$  , may be achieved by the use of some  $b_j$  , by the lemma, we have the  $2^{n+1}$  equations

$$M \begin{bmatrix} \beta \\ \beta_0 \\ \vdots \\ \beta_{01\dots n} \end{bmatrix} = 0$$

where the coefficient matrix  $M$  has mutually orthogonal rows. (Indeed, this production of possible  $\pm$  patterns - setting  $a_0 = \pm 1, a_1 = \pm 1, \dots$  in the monic words built from  $1, a_0, a_1, \dots, a_n$ , is one of the standard ways of producing a Hadamard matrix of side  $2^{n+1}$ .) Thus, each of the  $\beta$ 's is zero, and  $A_i$  is the zero of  $F(A)$  for each  $i$ . Hence  $r$  is zero. The module is faithful.

(ii)  $V$  is an irreducible  $F(G)$ -module.

Proof. Take any element  $v = \sum_{i=0}^n \alpha_i b_i$  in  $V$ , in which not all the  $\alpha_i$  are zero - say  $\alpha_k \neq 0$ . Then, by the lemma, and the symmetry of the array, there is an  $a_r$  such that

$$b_j a_r = \begin{cases} -b_k & \text{for } j = k, \\ +b_j & \text{for } 0 \leq j \leq n \text{ but } j \neq k. \end{cases}$$

So  $v\{e-a_r\} = 2\alpha_k b_k$ , so  $vF(G)$  contains  $b_k$ , and since by the action of  $\Sigma$ ,  $b_k F(G)$  contains all basic elements of  $V$ , we have  $vF(G) = V$  and hence  $V$  is irreducible.

It only remains to describe the algorithm to produce  $S$  and hence establish the lemma. We note first that if  $q = 2^k$ ,  $k > 0$ , and  $2^m > q$ , then the blocks along the top row of the array

$$(2^m, \dots, 2^{m+q-1}), (2^{m+q}, \dots, 2^{m+2q-1}), (2^{m+2q}, \dots), \dots$$

are "reflected", in the sense that the first  $q$  elements in the columns headed  $2^m; 2^{m+q}; 2^{m+2q}; \dots$  are just these blocks, in the same internal order. This follows since the permutations sending

0 to 1, 0 to 2, ..., 0 to  $2^{k-1}$  cannot change the binary digits past the  $2^{(k-1)}$ -digit. We start by flagging 2, 4, 5, 6, 7 as members of  $S$ : this deals with  $n = 0, 1$  and we have reached 7 as last flagged integer.

Assume we have flagged sufficient to justify the lemma for  $n = 0, 1, 2, \dots, 2^{k-1}-1$ , and that the last flag was placed at  $r$ . Set  $q = 2^k$ . (To start,  $k = 2, r = 7$ .)

(\*) Find the next power of 2, say  $2^m$ , such that  $2^m > r$  and  $2^m > q$ . Then allot flags within the next  $2^q$  blocks of length  $q$ ,

$$(2^m, 2^m+1, \dots, 2^m+q-1), (2^m+q, \dots, 2^m+2q-1), \dots$$

to produce one of each of the possible  $2^q$  patterns of flagging. This could be done in "dictionary" order, the first block totally unflagged, the last one totally flagged. We have now provided the  $\epsilon_n$  for

$n = 0, 1, 2, \dots, 2^k-1$ . Set  $r$  to the last integer flagged, set  $q = 2^{k+1}$ , and return to (\*). As noted in Section 3, this produces the set:

$$S = \{2, 5, 6, 7, 12, 17, 22, 27, 28, 29, \dots\},$$

and, in fact, the algorithm produces a plethora of columns of each required type for each  $n$ , and so the lemma is justified.

Consequently the irreducibility and fidelity of the module is established.

## Reference

- [1] Edward Formanek and Robert L. Snider, "Primitive group rings", *Proc. Amer. Math. Soc.* **36** (1972), 357-360.

Department of Mathematics,  
University of Newcastle,  
Newcastle, New South Wales.